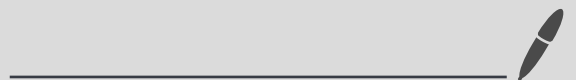


Theory of concurrency

2023/24

Lecture 7



# Approximate analysis of Petri nets using linear algebra

- A) invariants
- B) continuous reachability

hard problems:

- coverability
- boundedness
- finiteness

→ very hard problems:

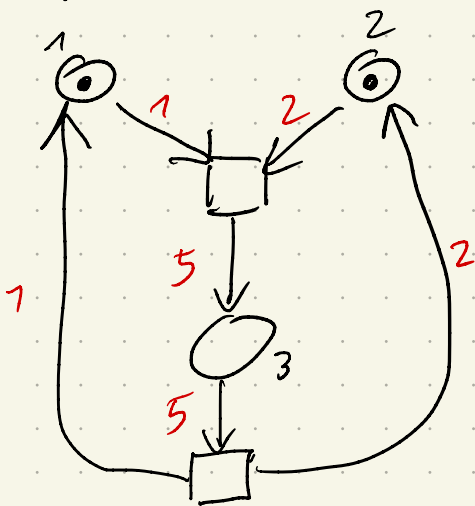
- reachability
- liveness

Model:

- general Petri nets
- counter automata w/o 0-tests
- VASS
- VAS = Petri nets without tight loops

A) invariants

Example:



$$\frac{1}{2} p_1 + \frac{1}{2} p_2 + p_3 = \text{const}$$

$$p_1 + p_3 = \text{const}$$

$$10 p_1 + 5 p_2 + 4 p_3 = \text{const}$$

$$18 p_1 + p_2 + 4 p_3 = \text{const}$$

$$5 p_1 + p_3 = \text{const}$$

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VAS :  $T \subseteq \mathbb{Z}^d$       $n := |T|$

$$N \in \mathbb{Z}[d \times n]$$

$$[x_1 \dots x_d] \cdot \begin{bmatrix} | & | & & | \\ t_1 & t_2 & \dots & t_n \\ | & | & & | \end{bmatrix} \cdot \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix}$$

Def : P-invariant :  $I \in \mathbb{Q}^d$  s.t.  $I \cdot N = \vec{0}$

T-invariant :  $I \in \mathbb{Q}^n$  s.t.  $N \cdot I = \vec{0}$

Petri net :

Fact :  $I : P \rightarrow \mathbb{Q}$  is a P-invariant  $\Leftrightarrow \forall t \in T$

$$\sum_{p \in t} I(p) \cdot F(p, t) = \sum_{p \in t} I(p) \cdot F(t, p)$$

$I : T \rightarrow \mathbb{Q}$  is a T-invariant  $\Leftrightarrow \forall p \in P$

$$\sum_{t \in p} I(t) \cdot F(t, p) = \sum_{t \in p} I(t) \cdot F(p, t)$$

3

Fact: Let  $V \subseteq \mathbb{Q}^d$  subspace generated by  $T$   
 $P$ -invariants = subspace  $\bar{V}$  orthogonal to  $V$

necessary  
 condition:

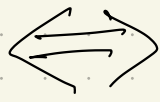
Fact:  $M \xrightarrow{*} M' \Rightarrow M \dashrightarrow^* M'$

$$M' = M + N \cdot X$$

for some  $X \in \mathbb{N}^n$   
 (NP-complete)

$\forall$   $P$ -invariant  $I$ .

$$I \cdot M = I \cdot M'$$



$$M' = M + N \cdot X$$

for some  $X \in \mathbb{Q}^n$

$$(M' - M) \text{ orthogonal to } \bar{V} \Leftrightarrow (M' - M) \in V$$

necessary condition:

all transitions are useful in  $M$

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Fact:

$M$  live  
no isolated places



$I \cdot M > 0$  for every semi-positive P-invariant  $I$   
( $I \geq \vec{0}, I \neq \vec{0}$ )

Proof: Let  $p \in P$  s.t.  $I(p) > 0$ .

Let  $M \xrightarrow{*} M'$  s.t.  $M'(p) > 0$

$$I \cdot M = I \cdot M' > 0 \quad \square$$

sufficient condition:

Fact: positive P-invariant  $\implies$  bounded

Proof: Suppose  $M \xrightarrow{*} M'$  Let  $p \in P$ .

$$I \cdot M = I \cdot M' \geq I(p) \cdot M(p) \quad \leftarrow \text{bounded by } \frac{I \cdot M}{I(p)} \quad \square$$

necessary and sufficient condition:

$M \xrightarrow{+} M \iff$  semi-positive T-invariant

B) continuous reachability

$$M \xrightarrow{*} M' \Rightarrow M \dashrightarrow^* M'$$

$$M' = M + N \cdot X$$

for some  $X \in \mathbb{N}^n$

(NP-complete)



(in P)

?

$$M' = M + N \cdot X$$

for some  $X \in \mathbb{Q}_+^n$



Def.

$$M \xrightarrow{t} M' \Leftrightarrow M' = M + t$$

$$M \dashrightarrow^{q \cdot t} M' \Leftrightarrow M' = M + q \cdot t$$

$$M \dashrightarrow^X M' \quad X \in \mathbb{Q}_+^n$$

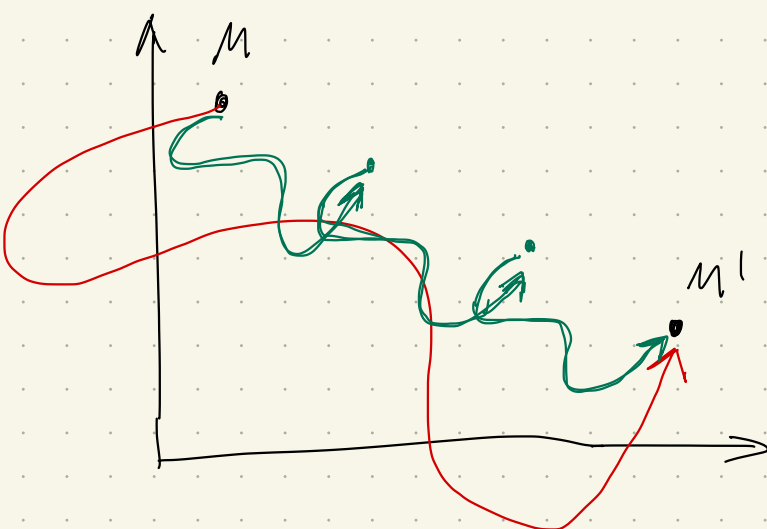
$t \in T$   
 $q \in \mathbb{Q}_+$   
 $M, M' \in \mathbb{Q}_+^d$

Lemma 1:  $M \dashrightarrow^X M' \Leftrightarrow M' = M + N \cdot X$

Proof:

if

$$(*) \left\{ \begin{array}{l} \forall t \in T, \forall p \in t, \\ X(t) > 0 \Rightarrow M(p) > 0 \\ \forall t \in T, \forall p \in t, \\ X(t) > 0 \Rightarrow M'(p) > 0 \end{array} \right.$$



# Lemma 2:



continuous rec



short continuous rec



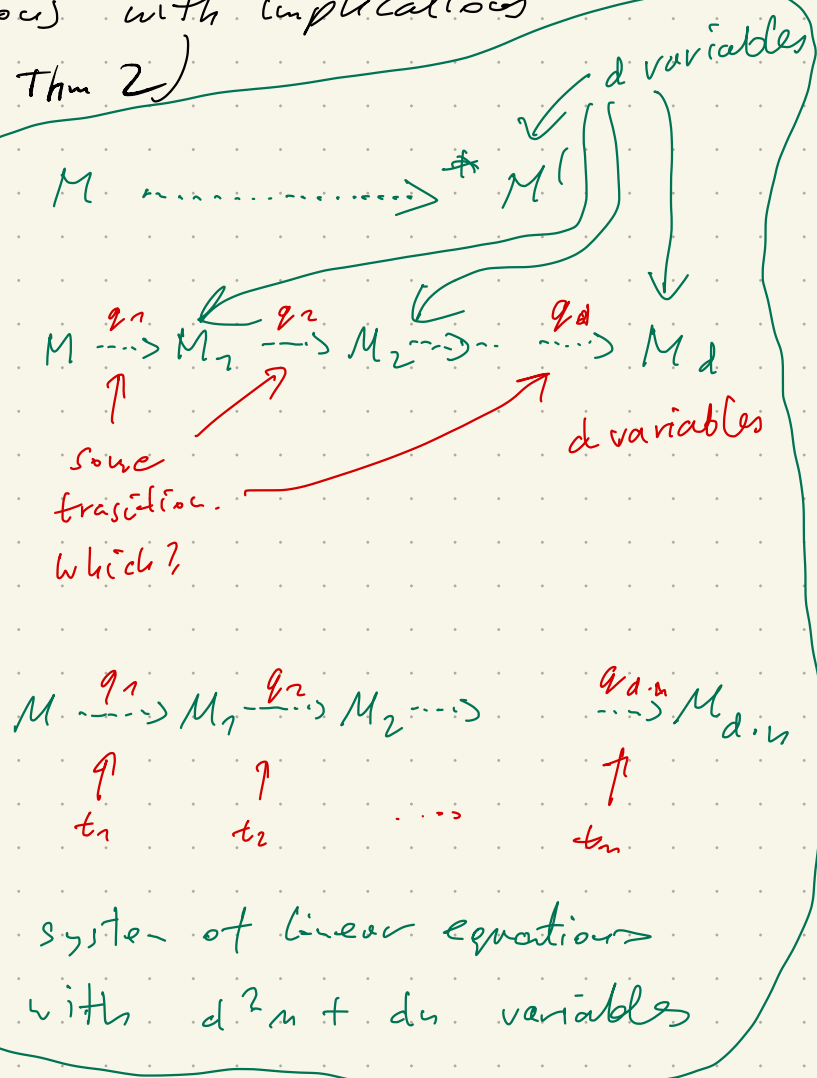
continuous rec satisfying (\*)

Theorem 1: Continuous reachability is in P.

Proof: encode in linear equations with implications  
(by Lemma 2 and Thm 2)

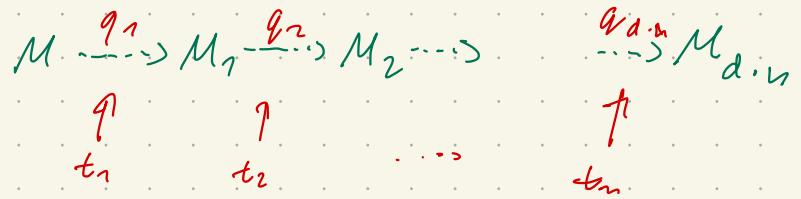
short continuous rec  
 $\leq d$  steps  
+

continuous rec  
satisfying (\*)



(by Lemma 1)

system of linear equations with implications  
with  $2d + n$  variables



system of linear equations  
with  $d^2 + d + n$  variables

(7)

Theorem 2: Solvability of linear equations  
with implications of the form

$$x_i > 0 \Rightarrow x_j > 0$$

over  $\mathbb{Q}$  is in  $\mathcal{P}$ .