

## Exam in information theory 8.02.2024. Theoretical part

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Please mark the answers that you consider right. Note that there may be more than one or none. You need not justify your choices. You **may optionally** add an argument for the point 5d (on the reverse side), which will be graded as *bonus*.

- Let  $w_n$ , for  $n \in \mathbb{N}$ , be a sequence of bit words, such that  $w_n \in \{0,1\}^n$ , and  $C_U(w_n) \geq n$ , for some fixed universal Turing machine  $U$ ; thus each  $w_n$  is Kolmogorov random w.r.t.  $U$ . Which of the properties below can hold for **infinitely many** words  $w_n$ ?
  - $w_n$  belongs to some Hamming code  $(2^m - 1, 2^m - m - 1)$ ,
  - $w_n$  belongs to some Huffman code (for some random variable),
  - $w_n$  contains a subword  $1^k$ , where  $k \geq \sqrt{n}$ ,
  - $w_n$  is a binary representation (possibly with leading zeros) of a number  $\lfloor \sqrt{k} \rfloor \cdot \lceil \log_2 m \rceil$ , for some integers  $k, m \geq 5$ .
- Consider random variables  $A, B, C$  with some joint distribution. Suppose that  $I(A; B|C) = I(A; B)$  and  $I(A; C|B) = I(A; C)$ . Which of the options below are consistent with this assumption (i.e., can happen)? Note that we ask about each option separately, not all of them together.
  - $I(B; C|A) \neq I(B; C)$ ,
  - $I(B; C|A) < I(B; C)$ ,
  - $I(B; C) = 0$ ,
  - $I(A; B|C) \neq I(A; C|B)$ .
- Suppose that, for a channel  $\Gamma$  with matrix  $\begin{pmatrix} p & q \\ q & p \end{pmatrix}$ , there exists a sequence of codes  $C_n \subseteq \{0,1\}^n$ , such that  $|C_n| \rightarrow \infty$ ,  $R(C_n) \rightarrow C_\Gamma$  and  $\Pr_E(\Delta_o, A_n) \rightarrow 0$  (for  $n \rightarrow \infty$ ), where  $A$  is a random variable taking values in  $C$  with uniform distribution, and  $\Delta_o$  is (attention!) the **ideal observer rule**. Then we can claim **for sure** that
  - $p = 0$  or  $p = 1$ ,
  - $p \neq \frac{1}{2}$ ,
  - $p > \frac{1}{2}$ ,
  - such a sequence cannot exist.
- A matrix of a channel  $\Gamma$  has dimension  $8 \times 8$ , but only 8 of its values is different from zero. Then the capacity  $C_\Gamma$ 
  - can be an arbitrary real number in the set  $[0, 3]$ ,
  - $0 < C_\Gamma < 3$ ,
  - can only assume a value from some **finite** set,
  - $C_\Gamma = 3$ .
- Three friends  $A, B, C$  decided to eat together **20 donuts**<sup>1</sup> but some randomness enters in their feast. Let the random variables  $A, B, C$  take values in the set  $\{0, 1, \dots, 20\}$  and represent how many donuts each person eats; we assume that **all donuts** will be eaten. Then we can claim for sure that
  - $H(A, B, C) \leq \log 231$ ,
  - $H(A, B, C) = H(A + B, B + C, C + A)$ ,
  - $H(A|B + C) \leq I(A; B|C)$ ,
  - a disjunction holds: some pair of variables is dependent (i.e.,  $I(A; B) + I(B; C) + I(A; C) > 0$ ) or  $H(A) = H(B) = H(C) = 0$ .

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<sup>1</sup>Donut (or doughnut, in Polish: *pączek*) is a bakery that in Poland is traditionally eaten (in great quantities) on Fat Thursday (Tłusty Czwartek), which happens just **today**.