

# Multi-Level Aggregation with Delays and Stochastic Arrivals

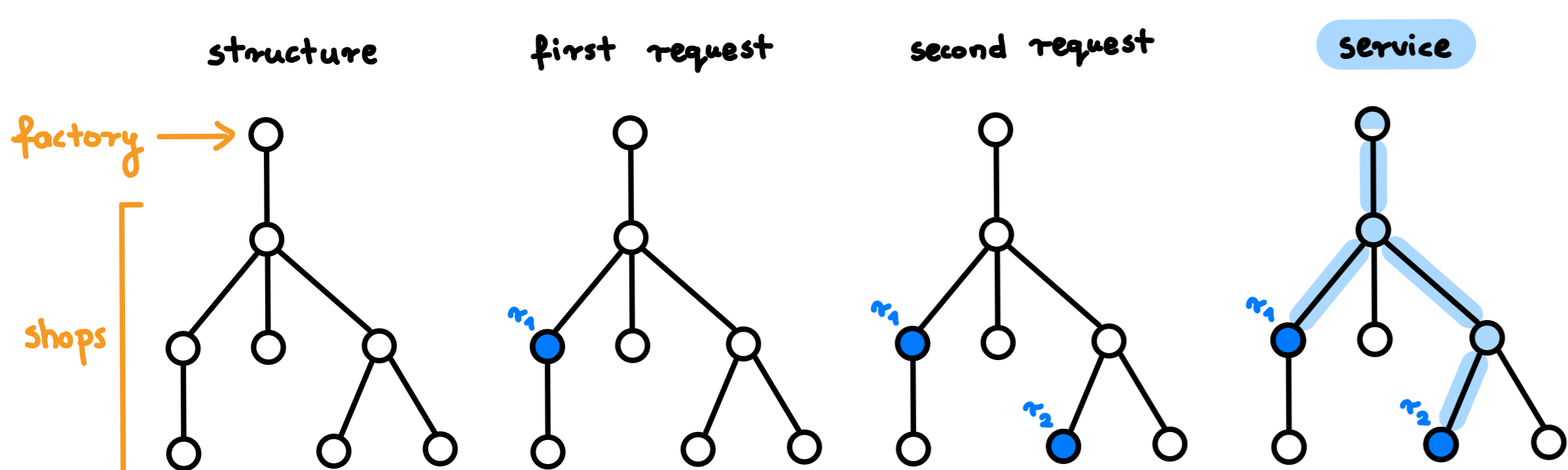
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## Motivation

Imagine a factory that needs to **schedule product deliveries** (services) to keep the cooperating shops running. Once a product is in **shortage** for a shop owner (agent), they inform the factory for **replenishment** (send a request). Each time the factory **schedules a service** to deliver some products, a **truck has to travel** from the factory to the locations of chosen agents. This incurs a cost proportional to the total travelling distance (we call it the service cost). Thus, to save on the delivery costs, it is beneficial to **accumulate** the replenishment **requests** from many stores and then deliver the ordered products altogether in one service. However, this **accumulated delay** in delivering the products may cause some agents **unsatisfied**. Typically, we measure this factor by looking at the **time gap between ordering and delivering** each requested product (and call it the delay cost of this request).

**Goal:** plan the delivery schedule in an online manner such that the total service cost and the total delay cost are minimized



## Problem statement

### Multi-Level Aggregation (MLA) with Delays [2]

- ▶ **edge-weighted** tree  $T$  rooted at  $\gamma$
- ▶  $V$  represents the set of  $T$ 's nodes,  $n := |V|$
- ▶ sequence  $\sigma$  of  $m$  **requests arriving online**
- ▶ request  $r$ : **location**  $\ell(r) \in V$ , **arrival time**  $t(r) \in \mathbb{R}^+$
- ▶ service  $s$  issued at time  $t$  to serve a set  $R$  of requests incurs the **delay cost** of  $\sum_{r \in R} (t - t(r))$  and the **service cost** being the weight of the minimum spanning tree containing all locations  $\ell(r)$  for  $r \in R$

**Target:** minimize the total cost produced by the online algorithm for serving all the arriving requests

## Previous results

- ▶ in **adversarial model**  $\rightarrow O(d^2)$ -competitive [1]
- ▶ **any algorithm** in this model  $\rightarrow \Omega(2 + \phi)$ -competitive [3]
- ▶ offline problem is NP-hard, 2-approximation exists

## Beyond worst-case

**Assumption:** the **waiting time** between any two consecutive requests arriving at any node  $u$  follows an **exponential distribution**  $\text{Exp}(\lambda_u)$  with parameter  $\lambda_u \geq 0$

- ▶ the factory needs to **minimize the expected cost** it produces when dealing with a random input sequence of requests arriving over some time horizon  $\tau$  for large  $\tau$
- ▶ to evaluate the performance of any algorithm  $A$  on stochastic input, we use the **ratio-of-expectations**  $\rightarrow$  the ratio of the expected costs of  $A$  and the optimal offline solution (OPT)

## Our main result

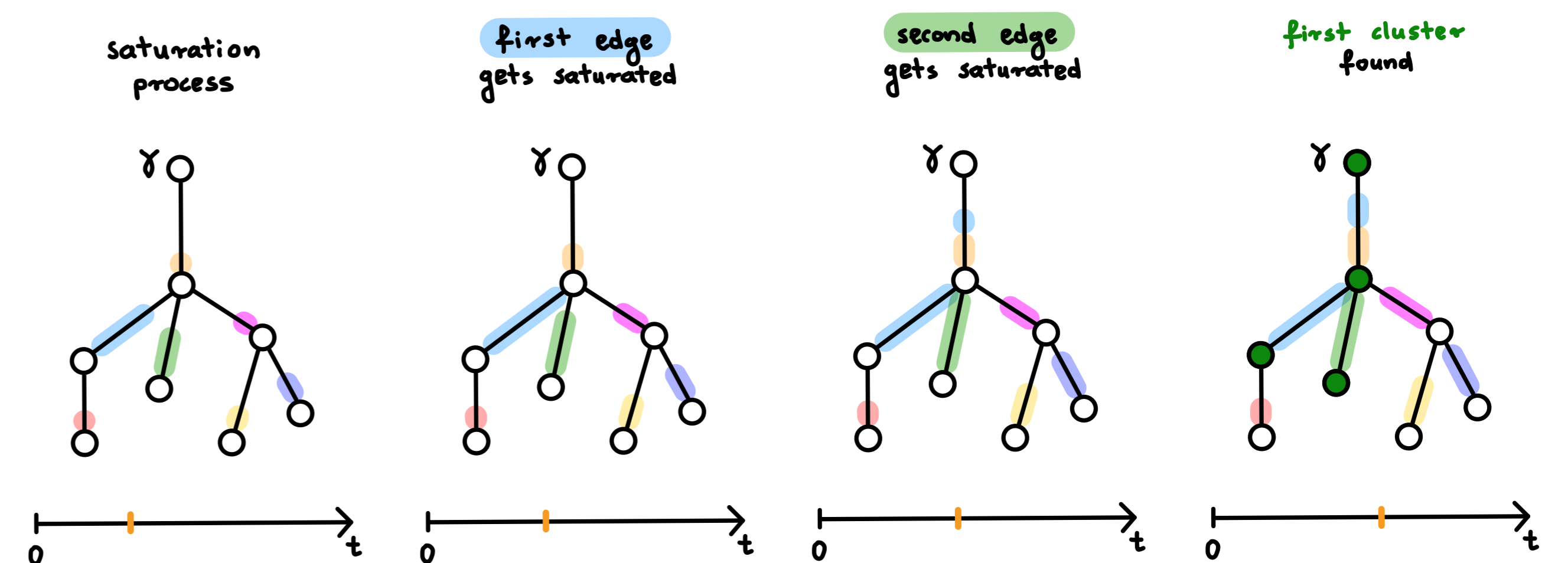
**Theorem.** For MLA with delays in the Poisson arrival model, there exists an algorithm with a **constant ratio-of-expectations**.

## Single-edge case

- ▶  $T$  has only **one edge**  $(u, \gamma)$  of weight  $w$
- ▶ denote the arrival rate of  $u$  by  $\lambda \rightarrow$  and consider **two strategies**
- ▶ **instant strategy:** serve each request as soon as it arrives
- ▶ efficient when the requests are **not so frequent**  $\rightarrow w \leq 1/\lambda$
- ▶ **periodic strategy:** group several requests arriving within a selected **period**  $p$  and **serve them together**
- ▶ efficient when requests **are frequent** enough  $\rightarrow w > 1/\lambda$
- ▶ we choose the period such that **the expected delay cost** generated by all the requests arriving within  $[0, p]$  **equals**  $w$ , i.e.,  $p = \sqrt{2w/\lambda}$

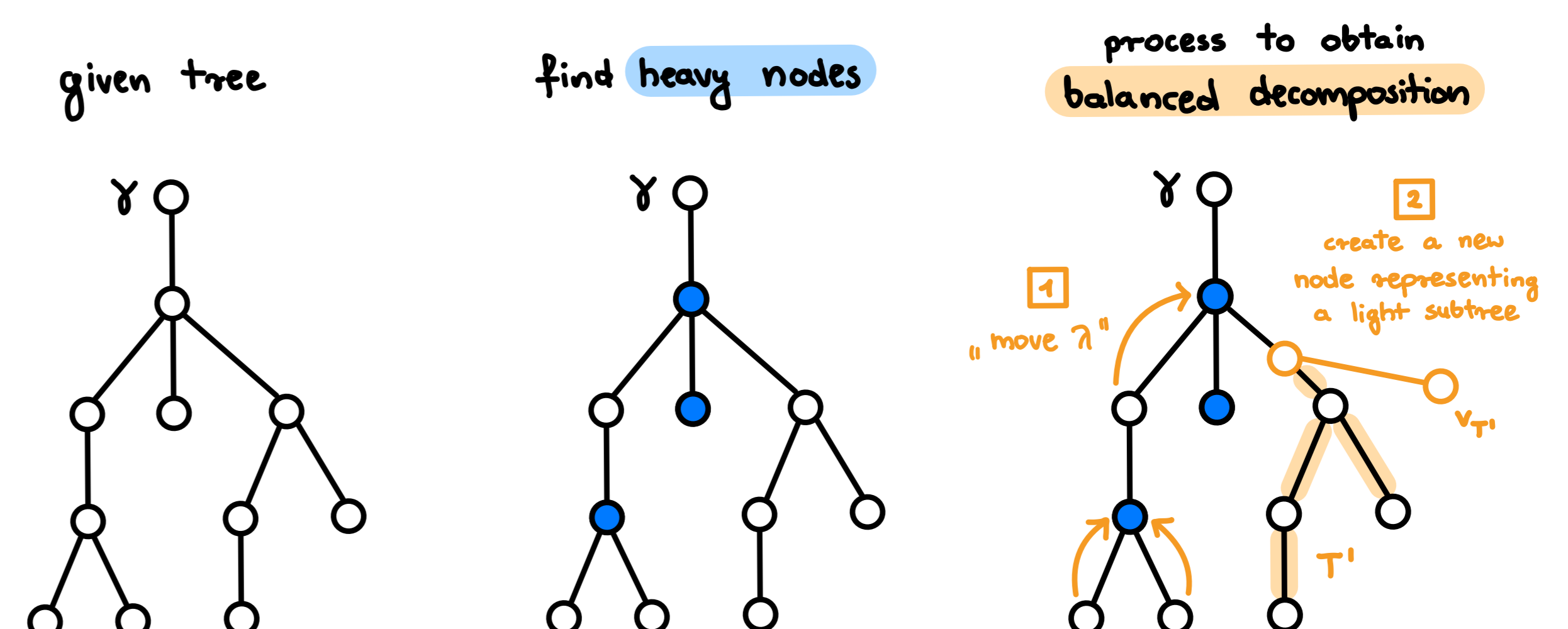
## Extending to more complex trees

- ▶ **instant strategy**  $\rightarrow$  trees for which the **average node-root distance**, weighted by the arrival rates, is **smaller** than the expected waiting time between two consecutive requests arrivals
- ▶ **periodic strategy**  $\rightarrow$  trees for which **each edge satisfies the single-edge case condition**; here, we use an edge-saturation-based process to assign each node its period



## General trees

- ▶ some of the nodes located **close to the root** should be served using **instant** strategy
- ▶ for the remaining ones, we need to **transform the tree** they form into a corresponding **heavy tree** and assign periods accordingly



## Bibliography

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