

# Piotr Glazar

**The Broadcast Storm Problem in a Mobile Ad Hoc Network**

# Introduction

- Wireless communication and mobile devices.
- Mobile Ad hoc NETWORK(MANET):
  - Unlimited possibility of movement.
  - No base station.
  - Multi-hop communication instead of single-hop.
- Used on battlefields or major disaster areas.

# Network

- Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA).
- No global network topology information.
- No collision detection in channel.
- Broadcast as the simplest way of sending data.
- Main issues:
  - Redundant broadcasts.
  - Heavy contention.

# Ideas

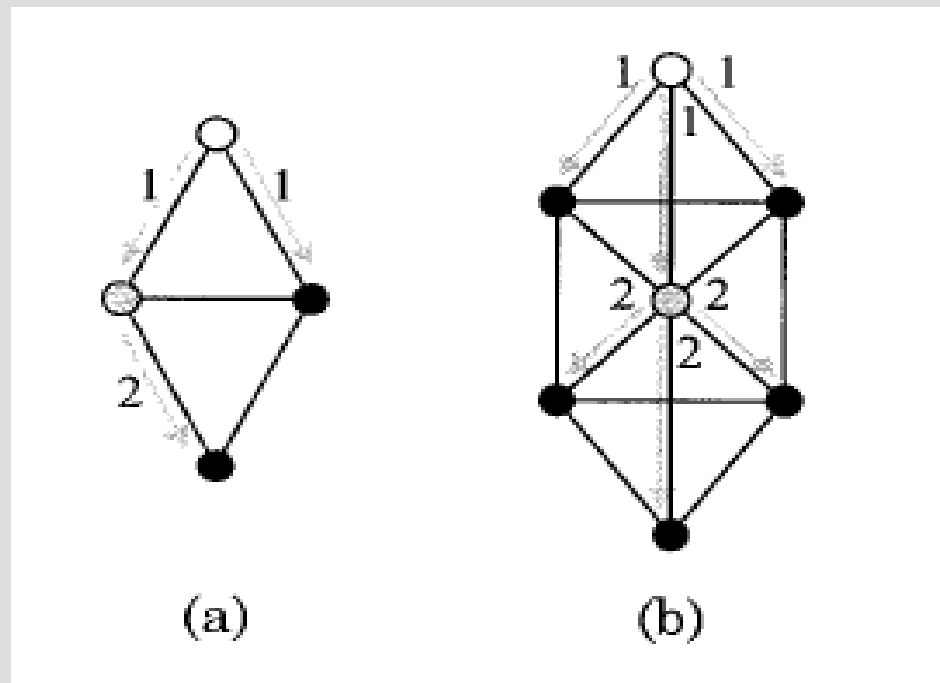
- Probabilistic scheme.
- Counter-based scheme.
- Distance-based scheme.
- Location-based scheme.
- Cluster-based scheme.

# Broadcast features

- Spontaneous,
- Unreliable,
- Messages can be distinguished from each other (source id and sequence number).

# Flooding

- Cost is  $n$  transmissions for  $n$  hosts in network.
- Many redundant broadcasts.
- Contention.



# Mathematics

- Rebroadcast can increase coverage: 61% is maximum, 41% on average.
- With third node we have 19% on average. For  $k > 3$  coverage is below 5%.
- Contention – 59% for three nodes, above 80% for at least 7.

# Probabilistic scheme

- On receiving a broadcast message for the first time, a host will rebroadcast it with probability  $P$ .
- For  $P = 1$  this solution is equivalent to flooding.
- Before rebroadcasting a node waits for random amount of time.



# Counter-based scheme

- Counter  $c$  tells us how many times the message was received.
- When  $c$  is greater than a fixed value a node doesn't send the message.
- 1. Initialize  $c = 1$  when the message is received for the first time.
- 2. Wait for a random amount of time and then send.
- 3. Message is on the air, exit.
- 4.  $c++$ , if  $c < C$  goto 2.
- 5. Don't send, exit.

# Distance-based scheme

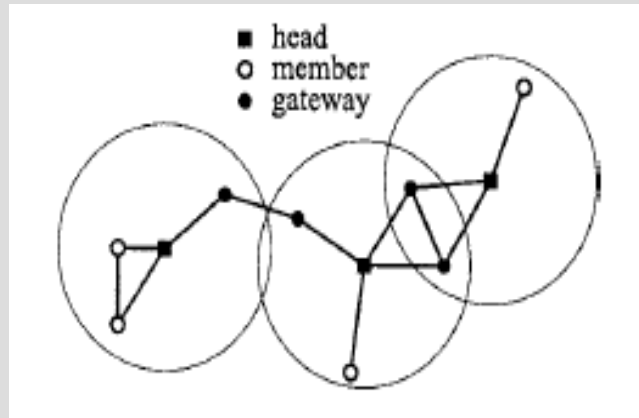
- Relative distance between nodes is used whether to send or not to sent the message.
- 1. Let  $d$  = distance to the broadcasting host when receiving the message for the first time.
- 2, 3 i 5 as before.
- 4.  $d = \min(d, \text{distance to the broadcasting host})$ , if  $d > D$  goto 2.
- Distance can be estimated from the signal strength.

# Location-based scheme

- We use additional information provided by GPS module. This allows us to count the additional coverage.
- Threshold  $0 < A < 0.61$ .
- Calculations are difficult, so we use convex polygons.
- Send the message if and only if a node is not located within a convex polygon formed by circles' centres.
- We can lose no more than 22% of coverage.

# Cluster-based solution

- Different kinds of nodes:



- Every node has its id which allows to determine role.
- Uses one of previous solutions whether to send or not to send.

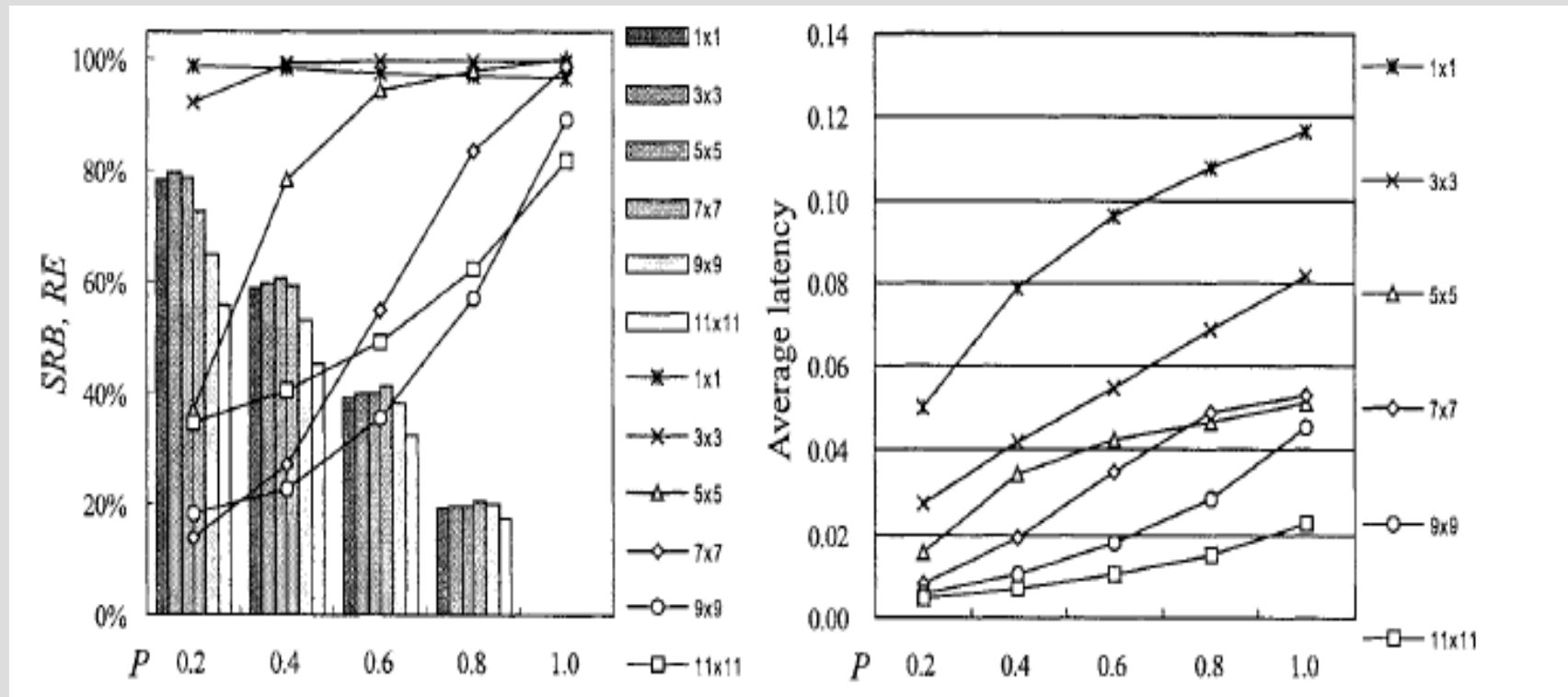
# Tests

- Transmission range – 500 metres.
- Message size – 280 bytes.
- Speed – 1Mbps.
- 100 nodes.
- Map size – squares: 1, 3, 5, 7, 9, 11 (x 500m).
- Nodes are placed in a random way, they can move around.
- Random waiting from 0 to 31 units.

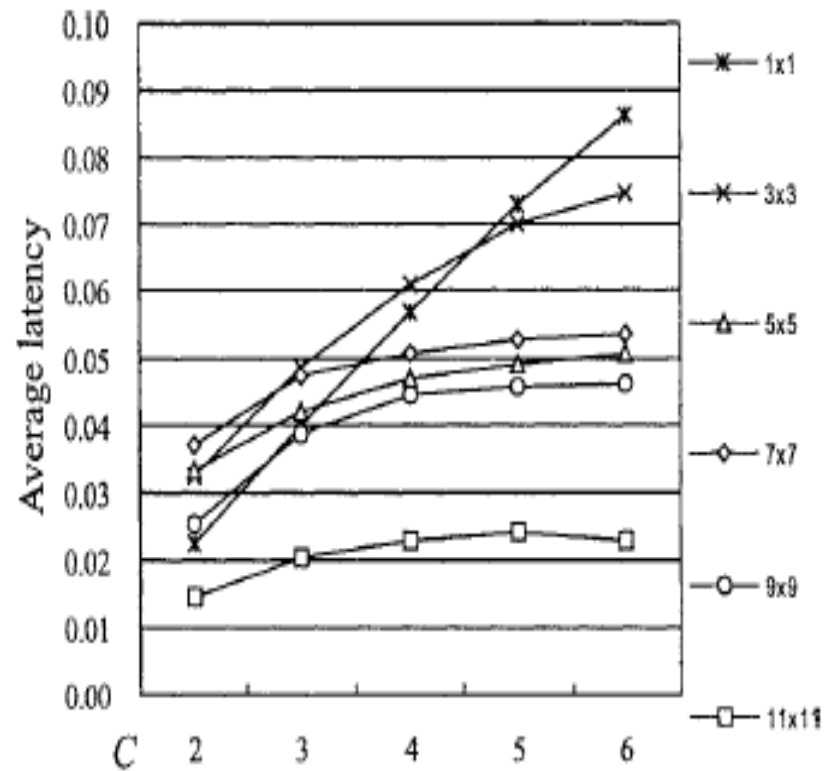
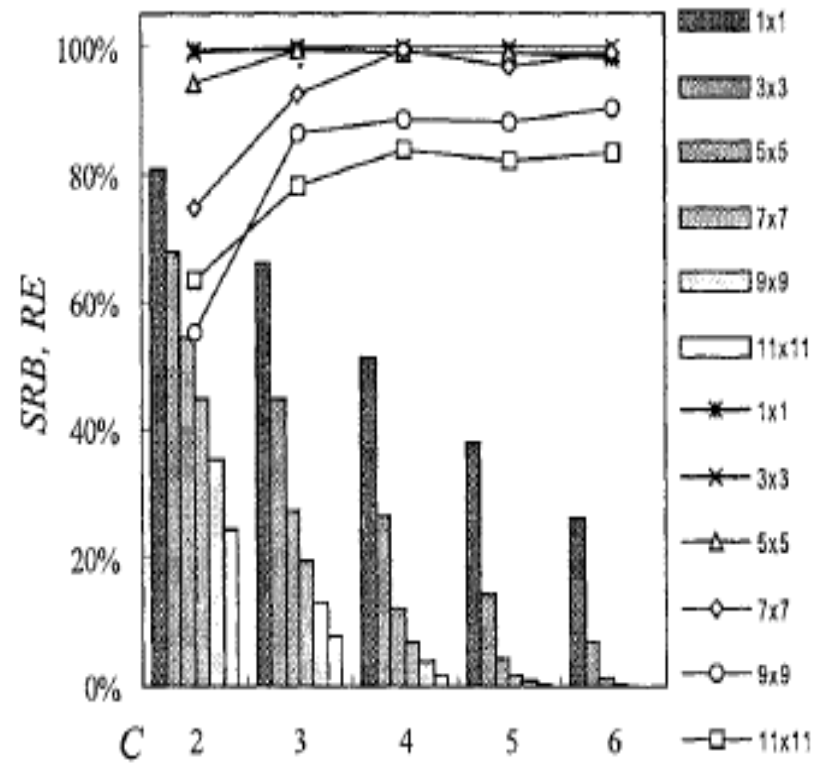
# Performance metrics

- Reachability (RE),
- Saved Rebroadcasts (SRB):  $(r - t) / r$ ,
- Average latency.

# Probabilistic scheme

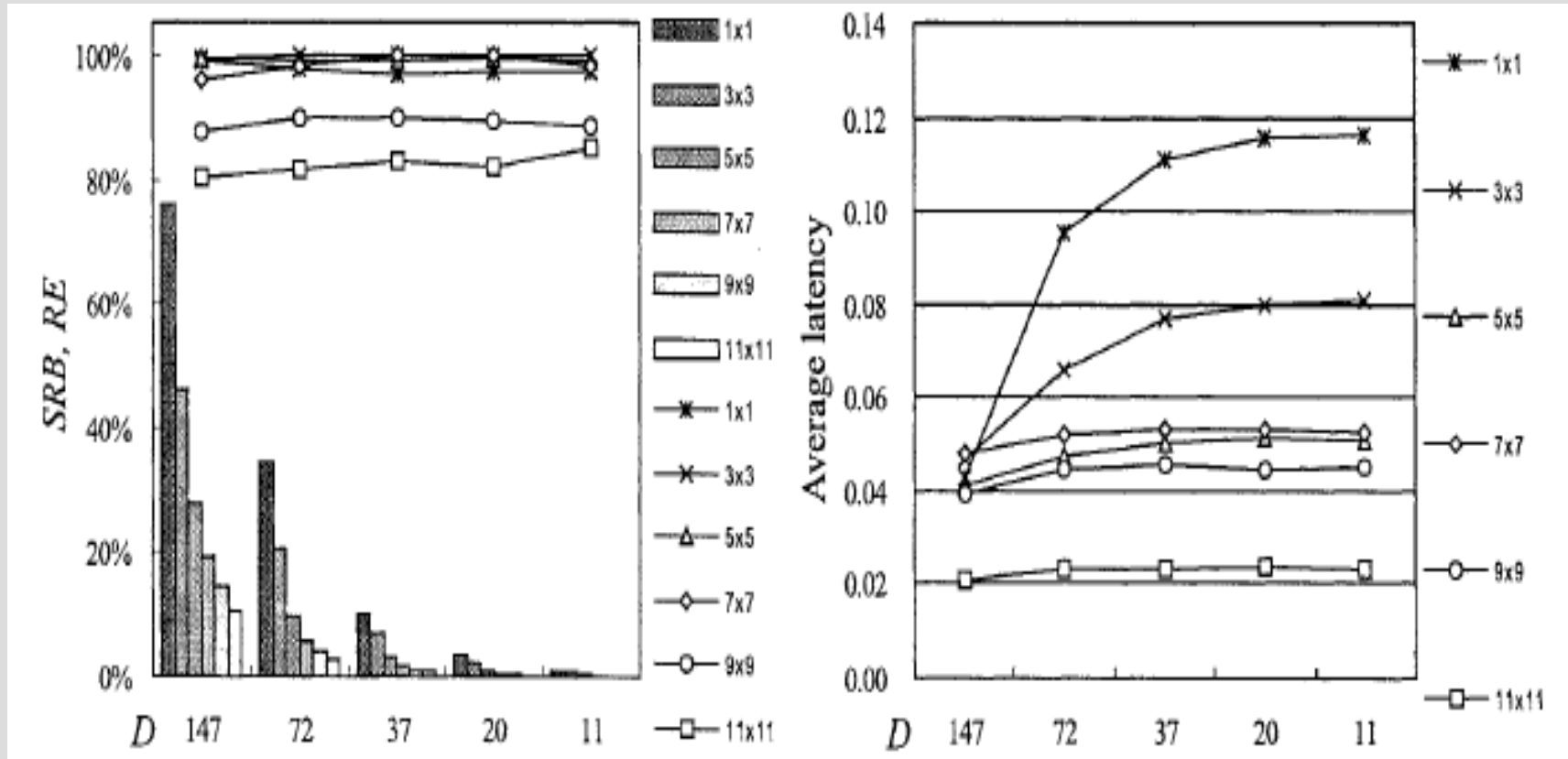


# Counter-based scheme

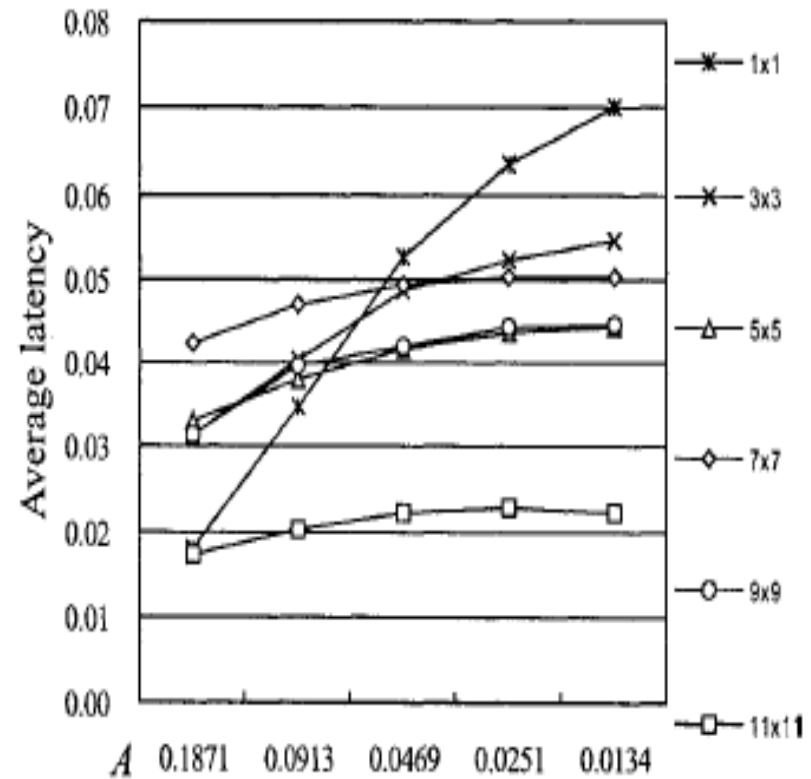
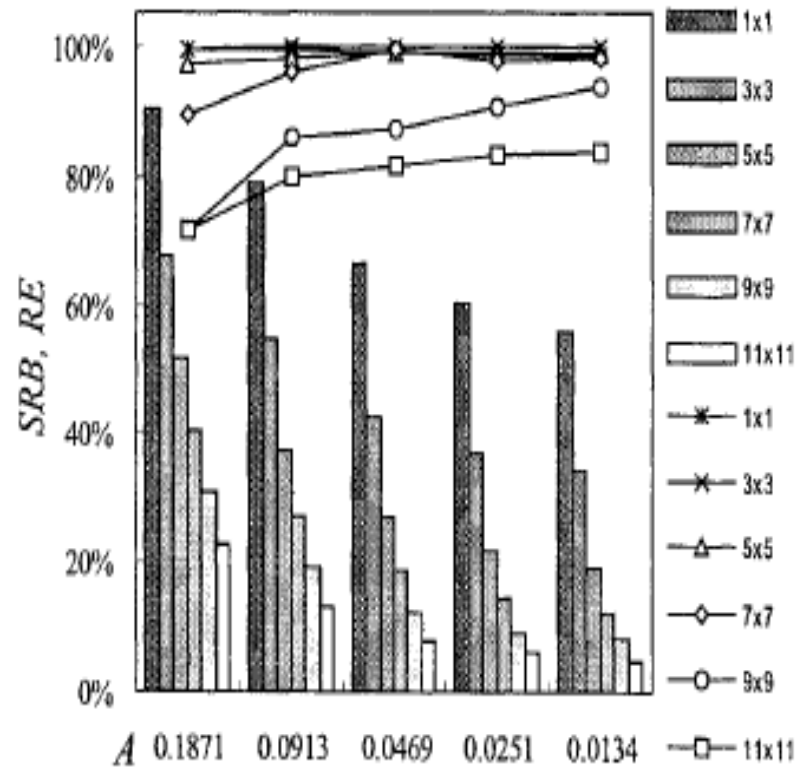




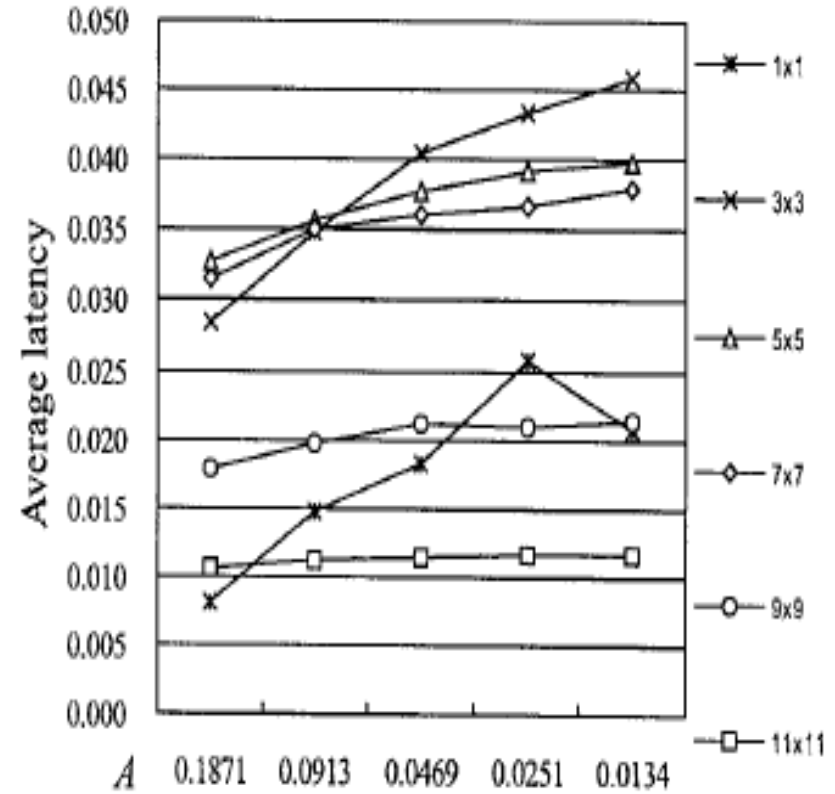
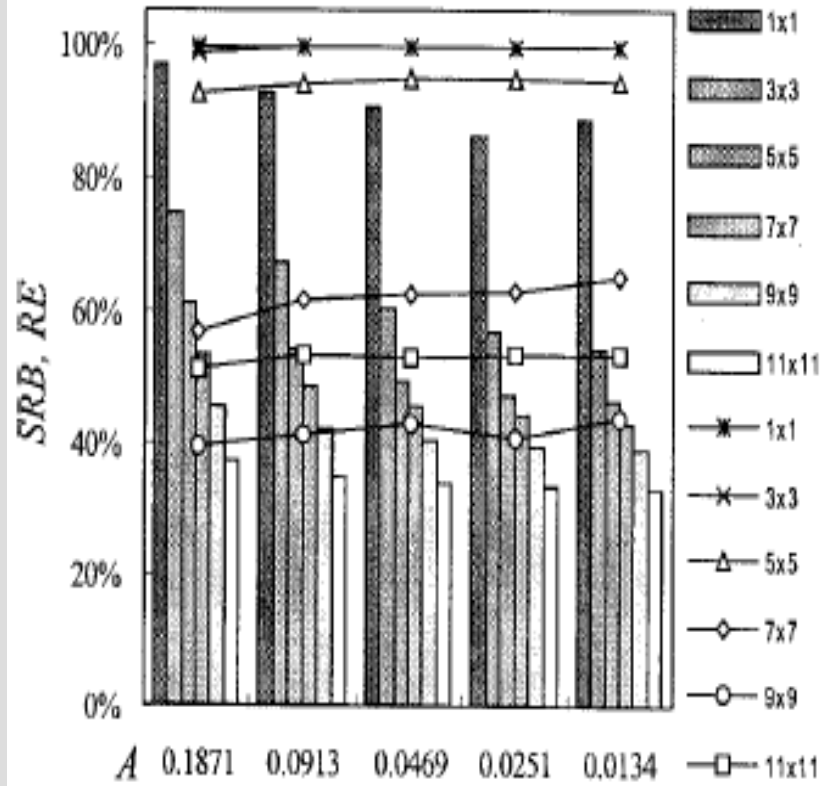
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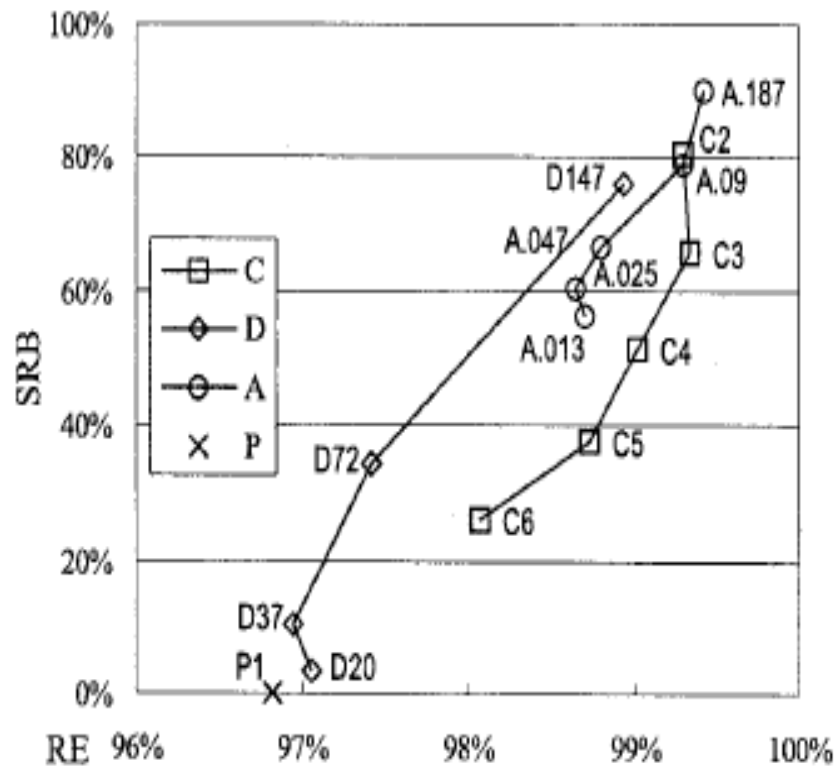
# Location-based scheme



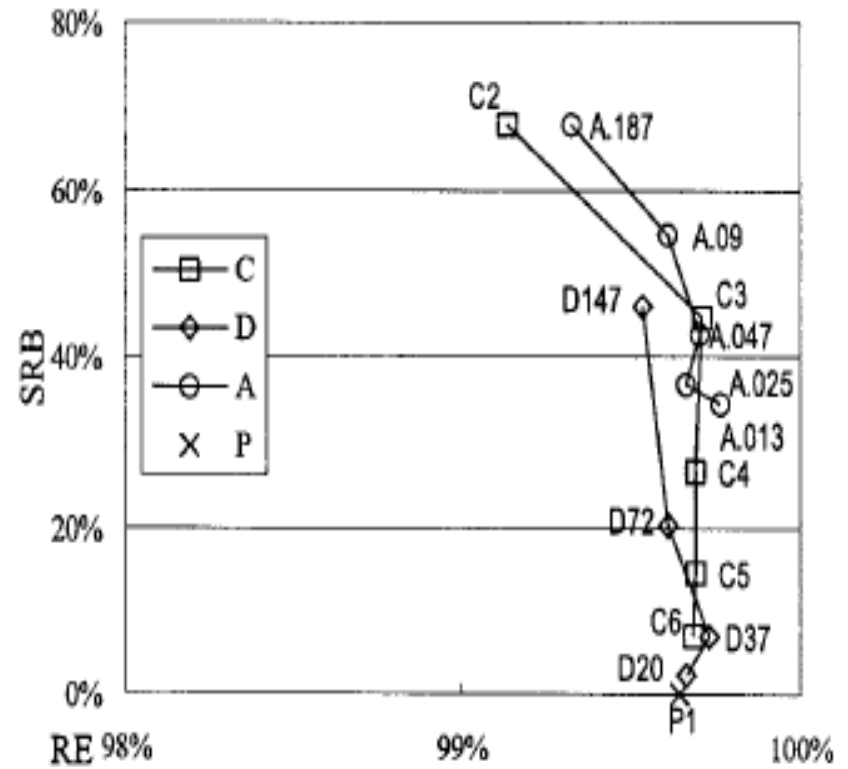
# Cluster-based



# RE / SRB

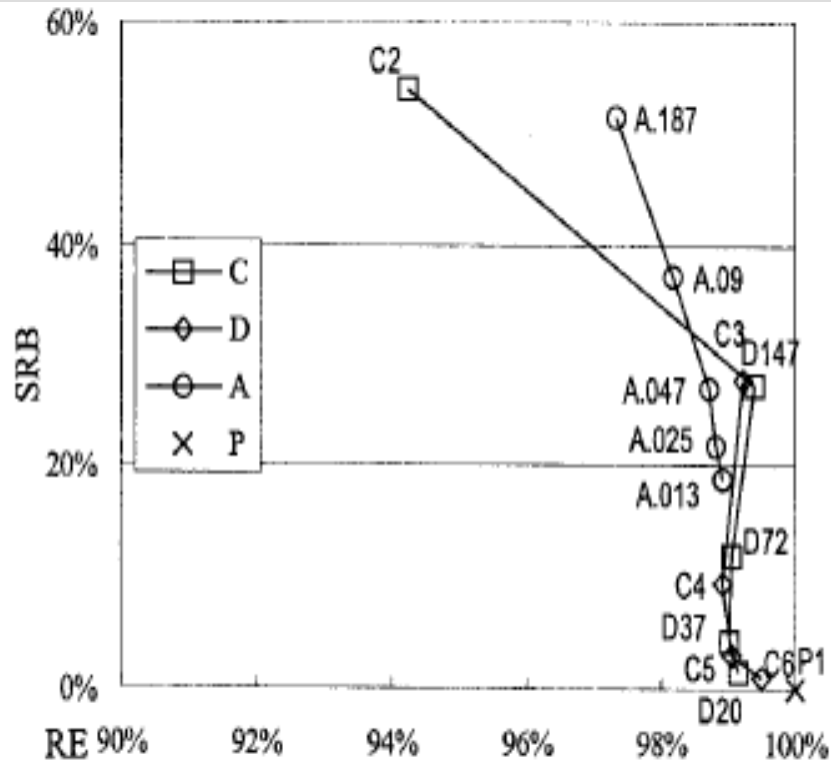


(a) 1x1 map

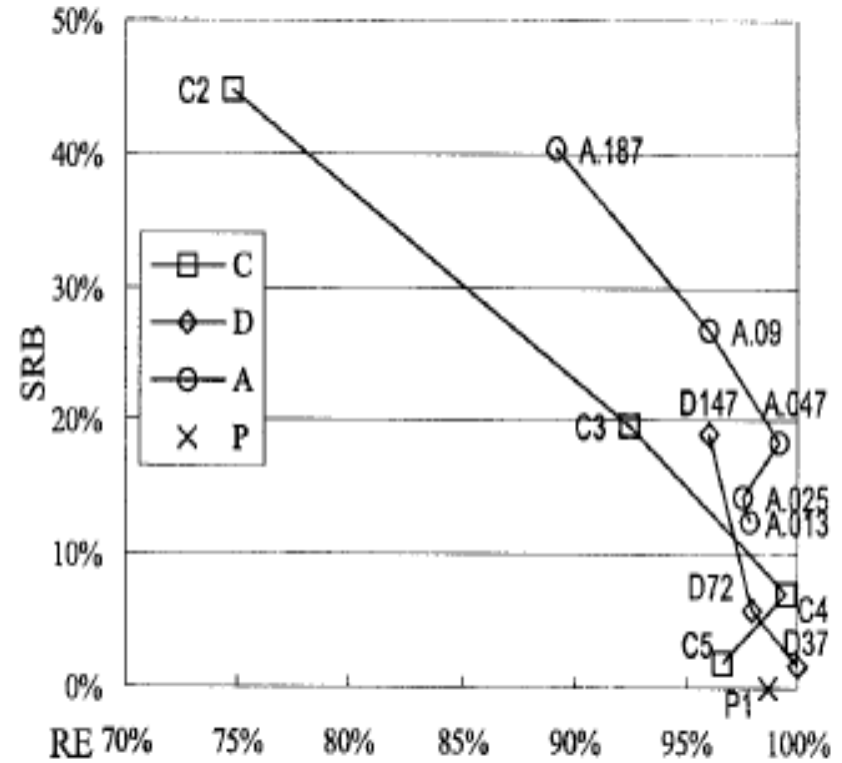


(b) 3x3 map

# RE / SRB

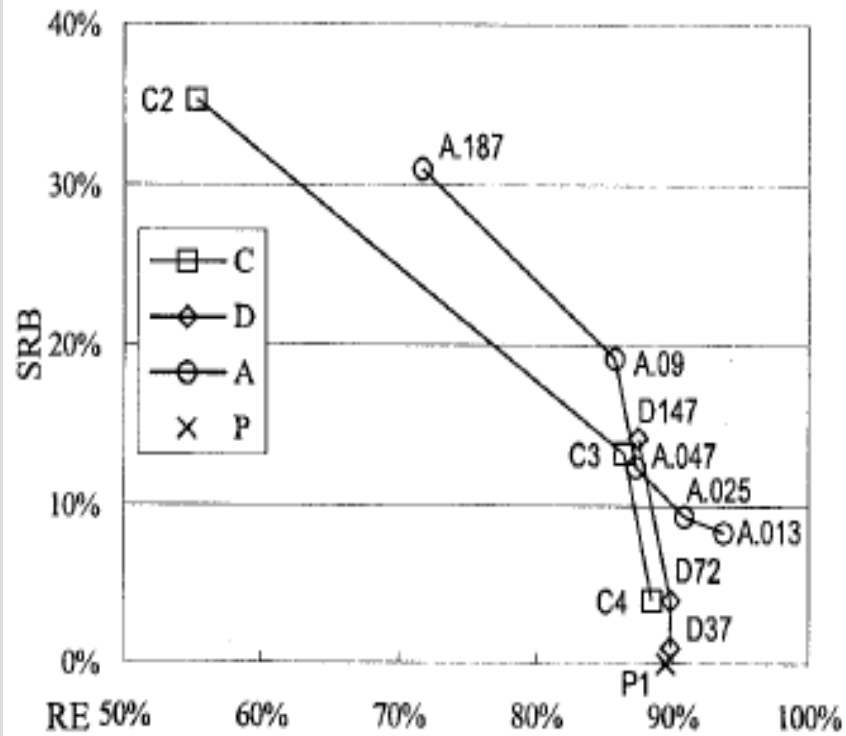


(c) 5x5 map

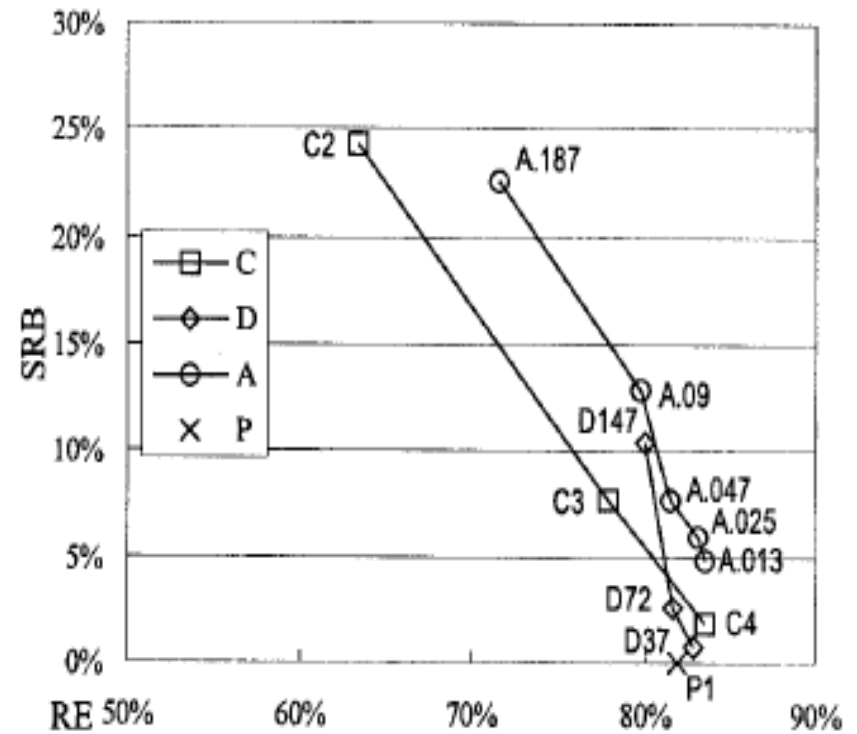


(d) 7x7 map

# RE / SRB



(e) 9x9 map



(f) 11x11 map

# Conclusion

- Only two solutions are worth mentioning.
- Counter-based scheme is easy to implement and provides satisfactory results.
- The best is location-based scheme.