

# Property A and duality in linear programming

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# Table of contents

1. What is ... property A?
2. Property A as a linear problem and the dual problem
3. Examples: primal and dual solutions
4. Primal relaxation: Cheeger constant
5. Applications: hypercubes and graphs with large girth

What is ... property A?

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# Property A

Property A was introduced in 2000 and turns out to be of great importance in many areas of mathematics [1]. Perhaps the most striking example is the following implication that follows from results in [4].

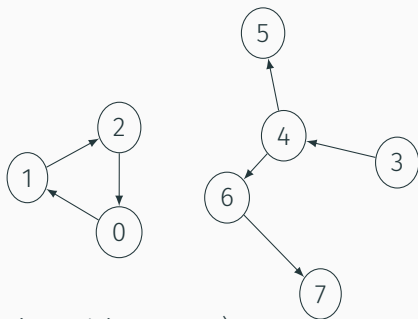
*If group  $G$  has Property A then the Novikov conjecture is true for all closed manifolds with fundamental group  $G$ .*

The Novikov conjecture asserts homotopy invariance of higher signatures of smooth manifolds. It is one of the most important unsolved problems in topology.

[1] P. Nowak, G. Yu, "What is ... property A?", Notices AMS (2008)

# Graph as a metric space

Path-length metric on a graph.



- Graphs are oriented ( $E \subset V \times V$ ).
- We allow infinite distance.
- If  $ij \in E$  is used to denote an edge, then  $i$  is the source and  $j$  is the target vertex of  $ij$ .

# Minimal variation of probability measures at scale $S$

**Optimization Problem I.** Let  $G = (V, E)$  be a graph and let  $S \geq 0$ . Find minimal  $\varepsilon = \varepsilon_{S,G}$  (variation) and a family  $\{\xi_i: V \rightarrow \mathbb{R}\}_{i \in V}$  of functionals (probability measures) such that

1. Each  $\xi_i$  is a probability measure, i.e.

$$\|\xi_i\|_1 = 1 \text{ and } \xi_i \geq 0 \text{ for each } i \in V;$$

2. Variation on edge  $ij$  does not exceed  $\varepsilon$ , i.e.

$$\|\xi_i - \xi_j\|_1 \leq \varepsilon \text{ for each } ij \in E;$$

3. Each  $\xi_i$  is supported by  $B(i, S)$ , i.e.

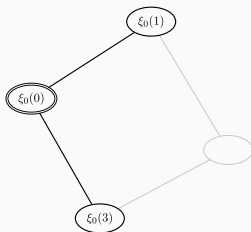
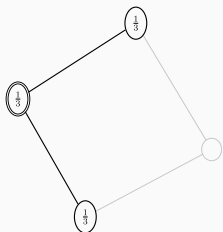
$$\text{supp } \xi_i = \{j \in V: \xi_i(j) > 0\} \subset B(i, S) \text{ for each } i \in V,$$

where  $B(i, S)$  is ball of radius  $S$  centered at  $i$ .

# Minimal example: square graph at scale 1

**Optimization Problem I.** Find minimal  $\varepsilon_{S,G} > 0$  with functionals  $\xi_i$  on  $G$  satisfying

1.  $\|\xi_i\|_1 = 1$  and  $\xi_i \geq 0$  for each  $i \in V$ ;
2.  $\|\xi_i - \xi_j\|_1 \leq \varepsilon$  for each  $ij \in E$ ;
3.  $\text{supp } \xi_i = \{j \in V : \xi_i(j) > 0\} \subset B(i, S)$  for each  $i \in V$ .



	$\xi_0$	$\xi_1$	$\xi_2$	$\xi_3$
0	$\frac{1}{3}$	$\frac{1}{3}$		$\frac{1}{3}$
1	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{1}{3}$	
2		$\frac{1}{3}$	$\frac{1}{3}$	$\frac{1}{3}$
3	$\frac{1}{3}$		$\frac{1}{3}$	$\frac{1}{3}$

This is optimal solution with objective  $\varepsilon = \frac{2}{3}$ .

Optimality of the solution is not trivial to show.

# Property A

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## Property A

Let  $G$  be a graph and for each  $S \geq 0$  let  $\varepsilon_{S,G}$  be the minimal variation of probability measures on  $G$  at scale  $S$ , i.e. the solution of Optimization Problem I at scale  $S$ . We say that  $G$  has **property A** iff

$$\lim_{S \rightarrow \infty} \varepsilon_{S,G} = 0.$$



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- To prove property A for  $G$  it is enough to find upper bounds  $\varepsilon_{S,G} \leq \hat{\varepsilon}_{S,G}$  such that

$$\lim_{S \rightarrow \infty} \hat{\varepsilon}_{S,G} = 0.$$

- To prove that  $G$  does not have property A we have to show that

$$\limsup_{S \rightarrow \infty} \varepsilon_{S,G} > 0.$$

so we have to consider optimal solutions  $\varepsilon_{G,S}$ .

# Reduction to finite graphs

- If  $S \geq \text{diam } G$ , then  $\varepsilon_{S,G} = 0$ , so property A is trivial for finite graphs.

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## Theorem

Let  $G$  be a graph and assume that  $G = \bigcup_{n \in \mathbb{N}} G_n$ , with  $G_1 \subset G_2 \subset G_3 \subset \dots$  an ascending sequence of convex subgraphs of  $G$ . Let  $\varepsilon_{S,G_n}$  be the minimal variation of probability measures at scale  $S$  for graph  $G_n$ . Graph  $G$  has property A iff

$$\lim_{S \rightarrow \infty} \lim_{n \rightarrow \infty} \varepsilon_{S,G_n} = 0.$$

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$$\lim_{S \rightarrow \infty} \lim_{n \rightarrow \infty} \varepsilon_{S,G_n} = 0.$$

Important case: disjoint sum of finite subgraphs.

# Graphs without property A

Graph  $G$  has property A iff

$$\lim_{S \rightarrow \infty} \lim_{n \rightarrow \infty} \varepsilon_{S, G_n} = 0.$$

To check that

$$\lim_{S \rightarrow \infty} \lim_{n \rightarrow \infty} \varepsilon_{S, G_n} \neq 0$$

we need a lower bound on  $\varepsilon_{S, G_n}$ .

## Property A as a linear problem and the dual problem

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# Linear programming formulation of minimal variance problem

**Optimization Problem I.** Find minimal  $\varepsilon_{S,G} > 0$  with functionals  $\xi_i$  on  $G$  satisfying

1.  $\|\xi_i\|_1 = 1$  and  $\xi_i \geq 0$  for each  $i \in V$ ;
2.  $\|\xi_i - \xi_j\|_1 \leq \varepsilon$  for each  $ij \in E$ ;
3.  $\text{supp } \xi_i = \{j \in V: \xi_i(j) > 0\} \subset B(i, S)$  for each  $i \in V$ .

**Primal problem.**

minimize  $e$

subject to  $\sum_{j \in V} x_{i,j} = 1$  for each  $i \in V$ ,

$$x_{i,j} = 0 \text{ for each } i \in V, j \in V \setminus B(i, S),$$

$$x_{j,k} - x_{i,k} \leq e_{ij,k} \text{ for each } ij \in E, k \in V,$$

$$x_{i,k} - x_{j,k} \leq e_{ij,k} \text{ for each } ij \in E, k \in V,$$

$$\sum_{k \in V} e_{ij,k} \leq e \text{ for each } ij \in E,$$

$$x_{i,j}, e_{ij,k}, e \geq 0$$

# The dual problem

Dual problem.

$$\text{maximize } \sum_{i \in V} \eta_i$$

subject to

$$\sum_{ij \in E} \kappa_{ij} \leq 1 ,$$

$$\varphi_{k,ij} \leq \kappa_{ij} \text{ for each } ij \in E, k \in V,$$

$$-\varphi_{k,ij} \leq \kappa_{ij} \text{ for each } ij \in E, k \in V,$$

$$\sum_{mi \in E, m \in V} \varphi_{k,mi} - \sum_{im \in E, m \in V} \varphi_{k,im} \geq \eta_i \text{ for each } k \in V, i \in B(k, S),$$

$$\eta_i, \varphi_{k,ij} \in \mathbb{R}, \kappa_{ij} \geq 0$$

## Primal

$$\min e$$

$$\text{s.t. } \sum_{j \in V} x_{i,j} = 1,$$

$$x_{i,j} = 0,$$

$$x_{j,k} - x_{i,k} \leq e_{ij,k},$$

$$x_{i,k} - x_{j,k} \leq e_{ij,k},$$

$$\sum_{k \in V} e_{ij,k} \leq e,$$

$$x_{i,j}, e_{ij,k}, e \geq 0$$

## Dual

$$\max \sum_{i \in V} \eta_i$$

s.t.

$$\sum_{ij \in E} \kappa_{ij} \leq 1,$$

$$\varphi_{k,ij} \leq \kappa_{ij},$$

$$-\varphi_{k,ij} \leq \kappa_{ij},$$

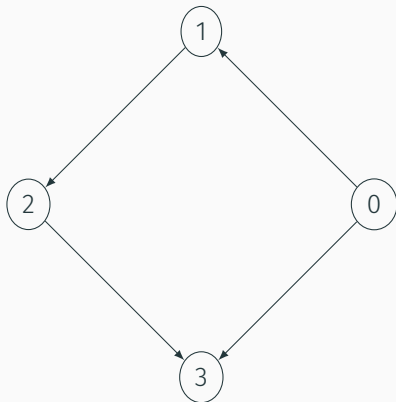
$$\sum_{mi \in E, m \in V} \varphi_{k,mi} - \sum_{im \in E, m \in V} \varphi_{k,im} \geq \eta_i,$$

$$\eta_i, \varphi_{k,ij} \in \mathbb{R}, \kappa_{ij} \geq 0$$

## Examples: primal and dual solutions

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# Square graph



## Square graph - proof of optimality

$$\|\xi_1 - \xi_0\|_1 = |\xi_1(0) - \xi_0(0)| + |\xi_1(1) - \xi_0(1)| + |\xi_1(2) - \xi_0(2)| + |\xi_1(3) - \xi_0(3)| \leq \varepsilon$$

↓

$$-\frac{1}{12}(\xi_1(0) - \xi_0(0)) + \frac{1}{12}(\xi_1(1) - \xi_0(1)) + \frac{1}{4}(\xi_1(2) - 0) - \frac{1}{4}(0 - \xi_0(3)) \leq \frac{1}{4}\varepsilon$$

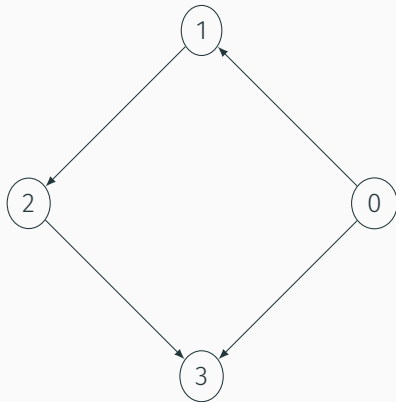
## Theorem

*Dual problem at scale  $S$  is dual to Primal problem at scale  $S$ . In particular, for each admissible solution of each problem, we have*

$$\sum_{i \in V} \eta_i \leq e$$

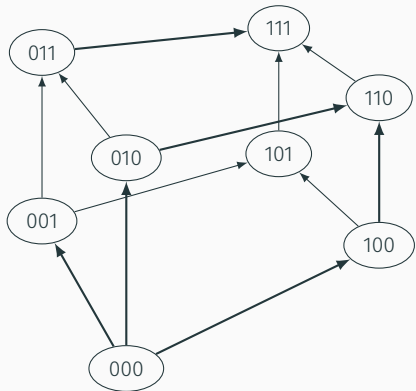
*and the optimal solutions are equal.*

# Square graph - solution by hand

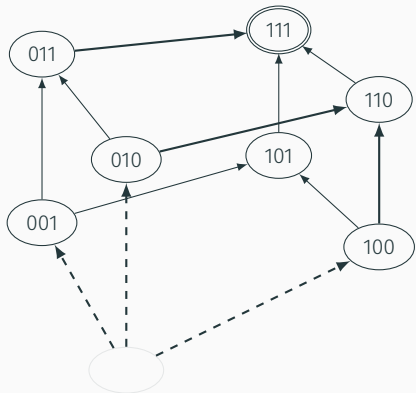




# Cube graph, $S = 2$ , primal solution



# Cube graph, $S = 2$ , dual solution



Primal relaxation: Cheeger  
constant

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# Capacity and supply is implicit if we know flows

Note that the capacity  $\kappa$  and supply  $\eta$  is implicit in the solution of the dual problem as the optimal values for chosen pseudo-flows  $\varphi_i$  can be easily computed.

## Dual Problem

$$\text{maximize } \sum_{i \in V} \eta_i$$

subject to

$$\sum_{ij \in E} \kappa_{ij} \leq 1 ,$$

$$\varphi_{k,ij} \leq \kappa_{ij} \text{ for each } ij \in E, k \in V,$$

$$-\varphi_{k,ij} \leq \kappa_{ij} \text{ for each } ij \in E, k \in V,$$

$$\sum_{mi \in E, m \in V} \varphi_{k,mi} - \sum_{im \in E, m \in V} \varphi_{k,im} \geq \eta_i \text{ for each } k \in V, i \in B(k, S),$$

$$\eta_i, \varphi_{k,ij} \in \mathbb{R}, \kappa_{ij} \geq 0$$

## Theorem

*Let  $G$  be a graph. Let  $\Gamma$  be a group that acts on  $G$  by automorphisms.*

*If  $\Gamma$  acts transitively both on edges and on vertices of  $G$ , then there exists an optimal solution of the dual problem such that  $\eta_i = \eta_j$  for each  $i, j \in V$  and  $\varepsilon_{ij} = \frac{1}{|E|}$  for each  $ij \in E$ .*

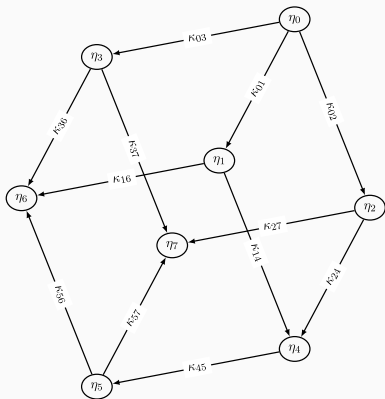
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This is not always the case.

# Forced equal capacities and supplies



$i$	$\sigma_{0,i}$	$\sigma_{1,i}$	$\sigma_{2,i}$	$\sigma_{3,i}$	$\sigma_{4,i}$	$\sigma_{5,i}$	$\sigma_{6,i}$	$\sigma_{7,i}$	$\min_k \sigma_{k,i}$
0	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$		$\frac{1}{28}$	$\frac{1}{28}$	$\eta_0 = \frac{1}{28}$
1	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$		$\eta_1 = \frac{1}{28}$
2	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$		$\frac{1}{28}$	$\eta_2 = \frac{1}{28}$
3	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$		$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\eta_3 = \frac{1}{28}$
4	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$		$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\eta_4 = \frac{1}{28}$
5		$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\eta_5 = \frac{1}{28}$
6	$\frac{1}{28}$	$\frac{1}{28}$		$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\eta_6 = \frac{1}{28}$
7	$\frac{1}{28}$		$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\frac{1}{28}$	$\eta_7 = \frac{1}{28}$
									$\sum = \frac{2}{7}$

$$\sigma_{k,i} = \sum_{m \in E, m \in V} \varphi_{k,mi} - \sum_{im \in E, m \in V} \varphi_{k,im}$$

# The dual problem with extra constraints

Uniform flows.

maximize  $|V| \cdot \eta$

subject to

$$\varepsilon_{ij,k} \leq \frac{1}{|E|} \text{ for each } ij \in E, k \in V,$$

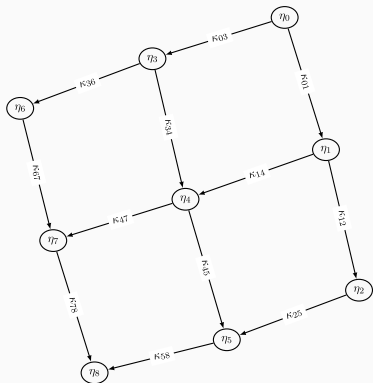
$$-\varepsilon_{ij,k} \leq \frac{1}{|E|} \text{ for each } ij \in E, k \in V,$$

$$\sum_{j \in V, ij \in E} \varepsilon_{ji,k} - \sum_{j \in V, ij \in E} \varepsilon_{ij,k} \geq \eta \text{ for each } k \in V, i \in B(k, S),$$

$$\eta, \varepsilon_{ij,k} \in \mathbb{R}$$



# $3 \times 3$ grid, Uniform Flows problem at scale $S = 1$



# Minimum isoperimetric number

Let  $S \subset V$ . The **edge boundary** of  $S$  is  $\partial S = E[S, V \setminus S]$ . For  $S \neq \emptyset$  we let

$$\varphi(S) = \frac{|\partial S|}{|S|}$$

be the **isoperimetric number** of  $S$ .

Let  $S \geq 0$  be a scale on a graph  $G$ . We let

$$\gamma(G, S) = \min_{T \subset B(i, S), i \in V, T \neq \emptyset} \varphi(T)$$

be the **Cheeger constant** of  $G$  at scale  $S$ .

## Theorem

*The optimal solution of Uniform Flow problem at scale  $S$  is equal to*

$$\frac{|V|}{|E|} \gamma(G, S),$$

*the Cheeger constant of  $G$  at scale  $S$  multiplied by  $\frac{|V|}{|E|}$ .*

## Minimal isoperimetric number

maximize  $\eta$

subject to

$$\varepsilon_{ij} \leq 1 \text{ for each } ij \in E,$$

$$-\varepsilon_{ij} \leq 1 \text{ for each } ij \in E,$$

$$\sum_{j \in V, ji \in E} \varepsilon_{ji} - \sum_{j \in V, ij \in E} \varepsilon_{ij} \geq \eta \text{ for each } i \in B(k, S),$$

$$\eta, \varepsilon_{ij} \in \mathbb{R}$$

## Minimal isoperimetric number - the dual

$$\text{minimize } \sum_{ij \in E} |a_i - a_j|$$

subject to

$$\sum_{i \in S} a_i = 1,$$

$$a_i = 0 \text{ for each } i \in V \setminus B(k, S),$$

$$a_i \geq 0 \text{ for each } i \in B(k, S)$$

### Theorem

*There exists an optimal solution of the above problem with all non-zero values equal.*

## Minimal isoperimetric number - the dual

$$\text{minimize } \sum_{ij \in E} |a_i - a_j|$$

subject to

$$\sum_{i \in S} a_i = 1,$$

$$a_i = 0 \text{ for each } i \in V \setminus B(k, S),$$

$$a_i \geq 0 \text{ for each } i \in B(k, S)$$

For such solution, if we take  $T = \{i: a_i > 0\}$ , then the value of the objective function is  $\frac{|\partial T|}{|T|} = \eta(T)$ . But this is the isoperimetric number of  $T$  - and it is the minimal one.

# Minimal isoperimetric number over $B(k, S)$

Therefore the minimal isoperimetric number dual is equivalent to:

## Minimal isoperimetric number - the dual reinterpreted

Maximize  $\eta$  such that for each  $T \subset B(k, S)$  we have

$$\eta \leq \frac{|\partial T|}{|T|}.$$

This is the Cheeger constant. Remember that we rescaled the original problem by  $\frac{|E|}{|V|}$ .

# The dual problem to Uniform flows

Uniform flows - Cheeger constant times  $\frac{|V|}{|E|}$ .

maximize  $|V| \cdot \eta$

subject to

$$\varepsilon_{ij,k} \leq \frac{1}{|E|} \text{ for each } ij \in E, k \in V,$$

$$-\varepsilon_{ij,k} \leq \frac{1}{|E|} \text{ for each } ij \in E, k \in V,$$

$$\sum_{j \in V, ji \in E} \varepsilon_{ji,k} - \sum_{j \in V, ij \in E} \varepsilon_{ij,k} \geq \eta \text{ for each } k \in V, i \in B(k, S),$$

$$\eta, \varepsilon_{ij,k} \in \mathbb{R}$$



# Mean property A - the dual problem to Uniform Flows

For each  $S > 0$  find  $\varepsilon_{S,G}$  and a family of functionals  $\{\psi_i\}$  on  $G$  satisfying

1.  $\{\psi_i\}$  has norm 1 *on average*, i.e.

$$\frac{1}{|V|} \left\| \sum_{i \in V} \psi_i \right\|_1 = 1,$$

and  $\psi_i \geq 0$  for each  $i \in V$ .

2.  $\{\psi_i\}$  has  $\varepsilon$ -variation *on average*, i.e.

$$\frac{1}{|E|} \sum_{ij \in E} \sum_{k \in V} |\psi_i(k) - \psi_j(k)| = \frac{1}{|E|} \sum_{ij \in E} \|\psi_i - \psi_j\|_1 \leq \varepsilon.$$

3.  $\text{supp } \psi_i \subset B(i, S)$  for each  $i \in V$ .

If  $\lim_{S \rightarrow \infty} \varepsilon_{S,G} = 0$ , then  $G$  has **mean property A**.

Applications: hypercubes and  
graphs with large girth

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## Theorem

Let  $Q_n$  be the  $n$ -dimensional hypercube graph. The minimal variation of probability measures for  $Q_n$  at scale  $S$  is

$$\varepsilon_{S, Q_n} = \frac{2 \binom{n-1}{S}}{\sum_{k=0}^S \binom{n}{k}}.$$

## Theorem

Let  $Q_n$  be the  $n$ -dimensional hypercube graph. The minimal variation of probability measures for  $Q_n$  at scale  $S$  is

$$\varepsilon_{S, Q_n} = \frac{2 \binom{n-1}{S}}{\sum_{k=0}^S \binom{n}{k}}.$$

For  $n = 3, S = 2$  we have

$$\varepsilon = \frac{2 \cdot \binom{2}{2}}{\binom{3}{0} + \binom{3}{1} + \binom{3}{2}} = \frac{2}{7}.$$

(We found this number before.)

**Corollary (P. Nowak, [2])**

*The disjoint union*

$$\coprod_{n \in \mathbb{N}} \{0, 1\}^n$$

*with  $\ell_1$  metric does not have property A.*

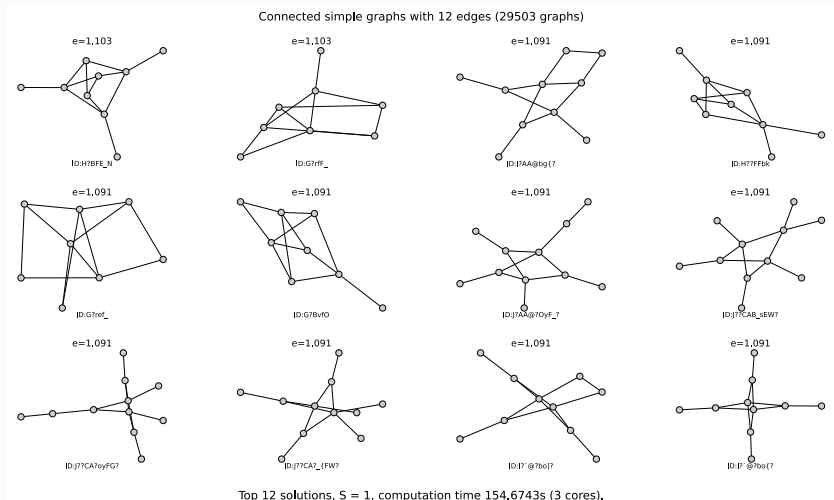
**Proof.**

The proof follows from the observation that for each  $S \geq 0$  we have

$$\lim_{n \rightarrow \infty} \frac{2^{\binom{n-1}{S}}}{\sum_{k=0}^S \binom{n}{k}} = 2.$$

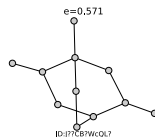
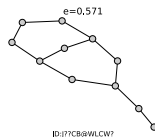
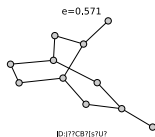
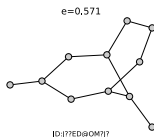
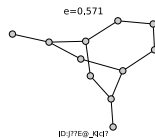
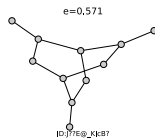
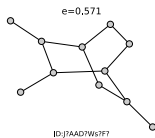
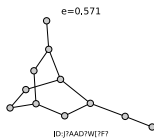
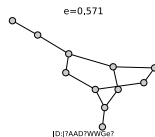
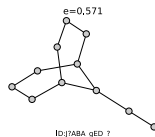
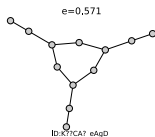
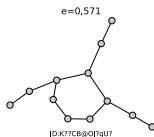


# Some experimental results and a quiz



# Some experimental results and a quiz

Connected simple graphs with 12 edges (29503 graphs)



Top 12 solutions,  $S = 2$ , computation time 207.7706s (3 cores).

# Large girth

## Theorem

Let  $G(d, c)$  be a  $d$ -regular graph with girth  $c$ . Let  $2S + 1 < c$ . The minimal variation of probability measures for  $G(d, c)$  at scale  $S$  is

$$\frac{2(d-1)^S(2-d)}{2-d(d-1)^S}.$$

## Corollary (R. Willett, [3])

Suppose  $d_i$  is a bounded sequence of integers with  $d_i \geq 3$  and suppose  $c_i$  is a sequence of integers going to infinity. Then, the disjoint union of the graphs  $G(d_i, c_i)$  fails to have property A.

## Proof.

The proof follows from the observation that

$$\lim_{S \rightarrow \infty} \frac{2(d-1)^S(2-d)}{2-d(d-1)^S} = 2 - \frac{4}{d}.$$





Thank you for your attention!



P. Nowak and G. Yu.

**What is . . . property A?**

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P. W. Nowak.

**Coarsely embeddable metric spaces without Property A.**

*J. Funct. Anal.*, 252(1):126–136, 2007.



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**Property A and graphs with large girth.**

*J. Topol. Anal.*, 3(3):377–384, 2011.



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**The coarse Baum-Connes conjecture for spaces which admit a uniform embedding into Hilbert space.**

*Invent. Math.*, 139(1):201–240, 2000.