Can algorithmics be useful in machine learning? Application of Banzhaf values to explain tree models

Piotr Sankowski

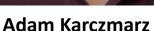






Main Result







Anish Mukherjee



Piotr Sankowski



Piotr Wygocki

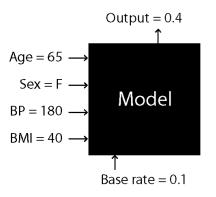


Tomasz Michalak

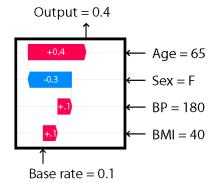
Karczmarz, Mukherjee, Wygocki, Michalak, Sankowski, "Improved Feature Importance Computation for Tree Models Based on the Banzhaf Value" The Conference on Uncertainty in Artificial Intelligence (UAI), Eindhoven, Netherlands, August 1-5, 2022.

SHAP and TreeSHAP



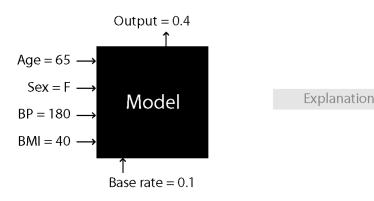


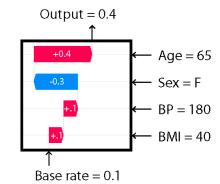
Explanation



SHAP and TreeSHAP







Lundberg et al. (2018, 2020) proposed TreeSHAP - an exact algorithm to compute the Shapley value-based explanations for tree models in $O(TLD^2 + n)$. Here:

- n is the number of features
- T the number of trees
- L the number of leaves
- D the maximum depth of a tree

Take away message

Technical contribution:

We advocate the Banzhaf value for tree models:

- 1. It can be computed noticeably faster
- 2. It seems to be more numerically stable
- 3. Our experimental comparison shows:
 - essentially the same global impacts
 - close explanations of individual predictions

Meta level:

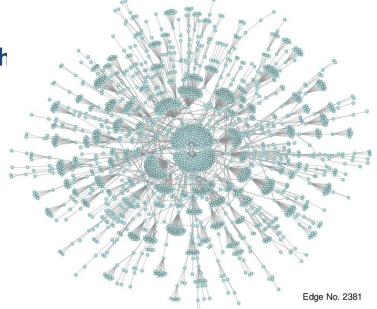
- Game theory and algorithmic view
- Interplay of the above areas with AI is growing
- Many more interesting problems to come

Scale-free Networks

Definition: An undirected graph G is called a power-law graph with parameter $\alpha>1$ if the fraction of vertices of degree k is proportional to $k^{-\alpha}$.

Theorem: If G is "power-law graph" then th heuristic finds maximum clique

- in polynomial time for for $\alpha > 3$,
- subexponential time for $2 < \alpha < 3$.



Pawel Brach, Marek Cygan, Jakub Lacki, Piotr Sankowski:

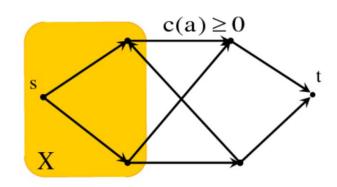
Algorithmic Complexity of Power Law Networks. SODA 2016: 1306-1325

Scale-free Networks

Definition: If Ω is a finte set, a function $f: 2^{\Omega} \rightarrow R$ is submodular when

• For every $S, T \subseteq \Omega$ we have that $f(S) + f(T) \ge f(S \cup T) + f(S \cap T)$.

Theorem: When a submodular function is decomposable into sum of simple submodular functions then its minimum can be found in time needed to solve the maximum flow problem.



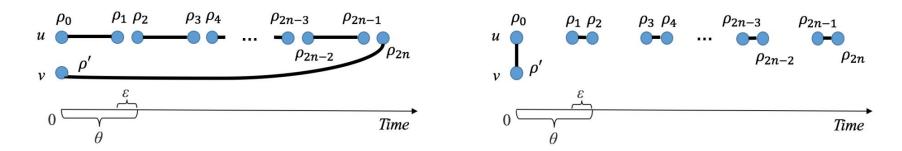
Kyriakos Axiotis, Adam Karczmarz, Anish Mukherjee, Piotr Sankowski, Adrian Vladu: Decomposable Submodular Function Minimization via Maximum Flow. ICML 2021: 446-456.

Li Chen, Rasmus Kyng, Yang P. Liu, Richard Peng, Maximilian Probst Gutenberg, Sushant Sachdeva: *Maximum Flow and Minimum-Cost Flow in Almost-Linear Time*. FOCS 2022: 612-623

Stochastic Arrivials

Definition: In the Min-cost Perfect Matching with Delays (MPMD) problem we need to match online requests by paying:

- the connection cost,
- the waiting time cost.



Theorem: For stochastic arrivals the greedy heuristic is constant competetive in expectation.

Mathieu Mari, Michał Pawłowski, Runtian Ren and Piotr Sankowski: *Online matching with delays and stochastic arrival times, AAMAS 2023.*

1. Values in Cooperative Game Theory

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Coalitional Games

Given the **set of agents**:

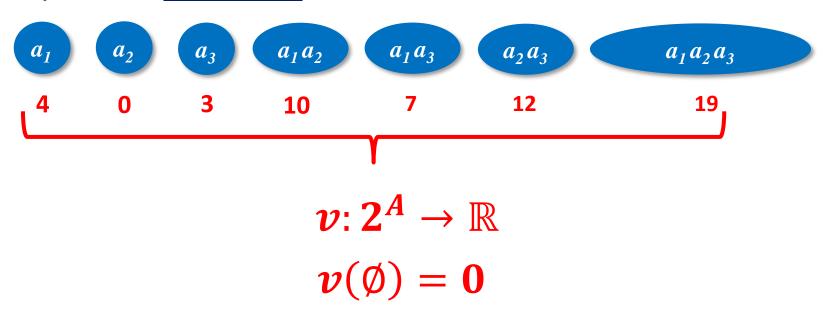
$$A = \{a_1, a_2, a_3\}$$

Coalitional Games

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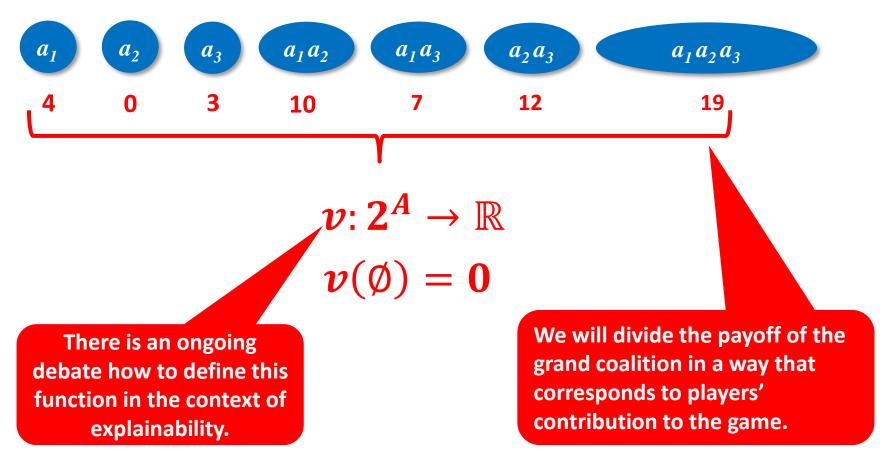


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Marginal contribution:

Let $C \subseteq A \setminus \{a_i\}$. Then:

$$MC(a_i, C) = v(C \cup \{a_i\}) - v(C).$$

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Marginal contribution:

Let $C \subseteq A \setminus \{a_i\}$. Then:

$$MC(a_i, C) = v(C \cup \{a_i\}) - v(C).$$

We are interested in a method that considers the marginal contributions of a player to all the coalitions in the game.

Shapley value

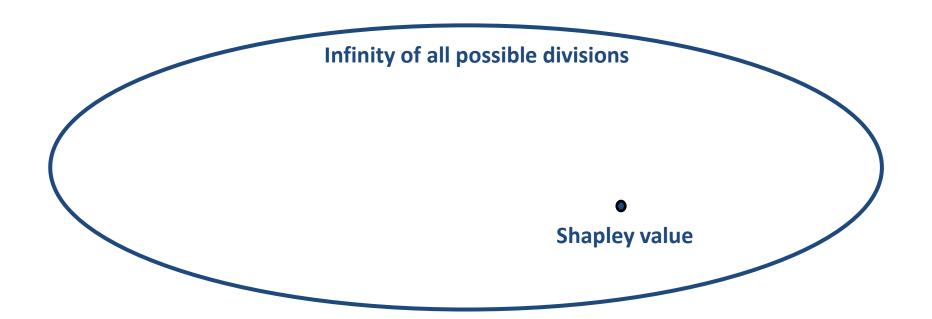
- Symmetry any two players who always contribute the same to all coalitions should get the same payoff (i.e. the same share of the grand coalition)
- Null player a player who does not contribute anything to any coalition should get nothing
- Additivity for additive games, the payoffs should be also additive
- Efficiency the <u>total value of the grand coalition</u> should be distributed among the players there should be no leftovers and we should not be able to distribute more than we have

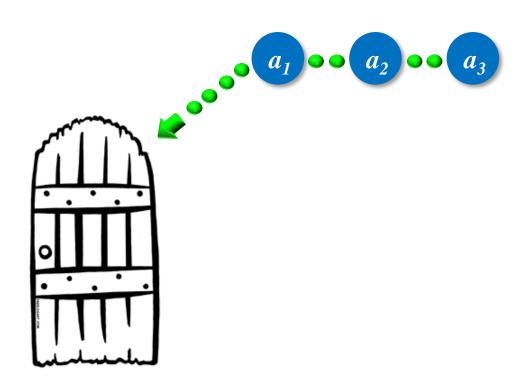
Shapley value

There exists unique value that satisfies Symmetry, Null player, Additivity and Efficiency. It is defined as follows:

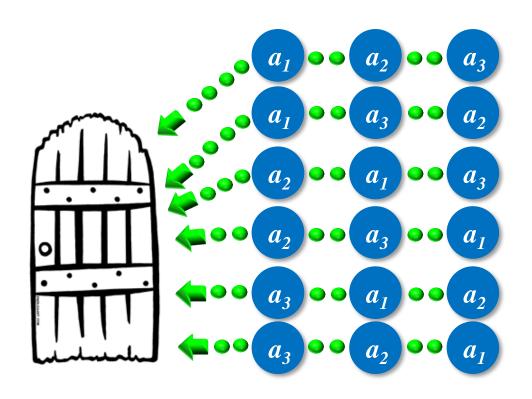
$$Sh_{i}(v) = \sum_{\mathbf{C} \subseteq A \setminus \{a_{i}\}} \frac{|C|! (|A| - |C| - 1)!}{|A|!} [v(C \cup \{a_{i}\}) - v(C)]$$

Taxonomy of Solutions

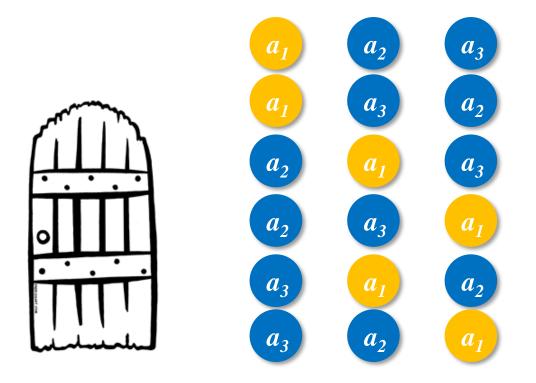




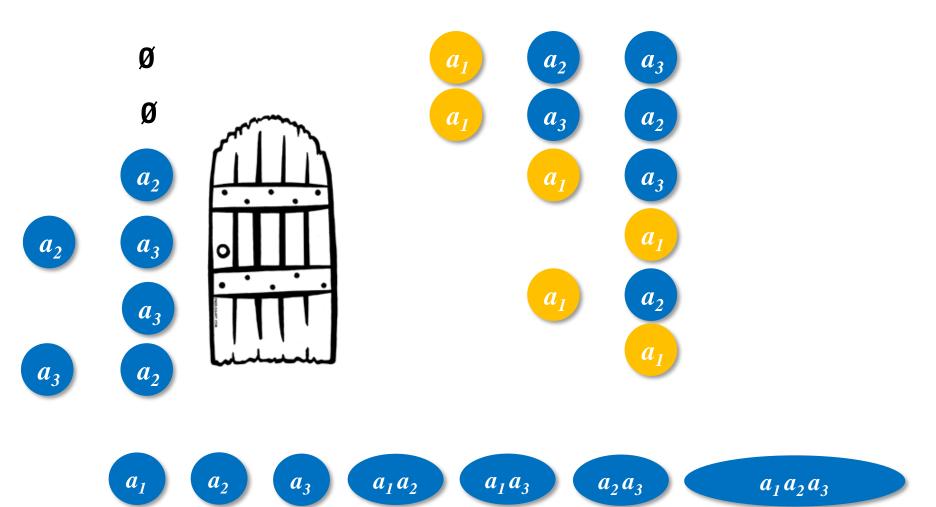


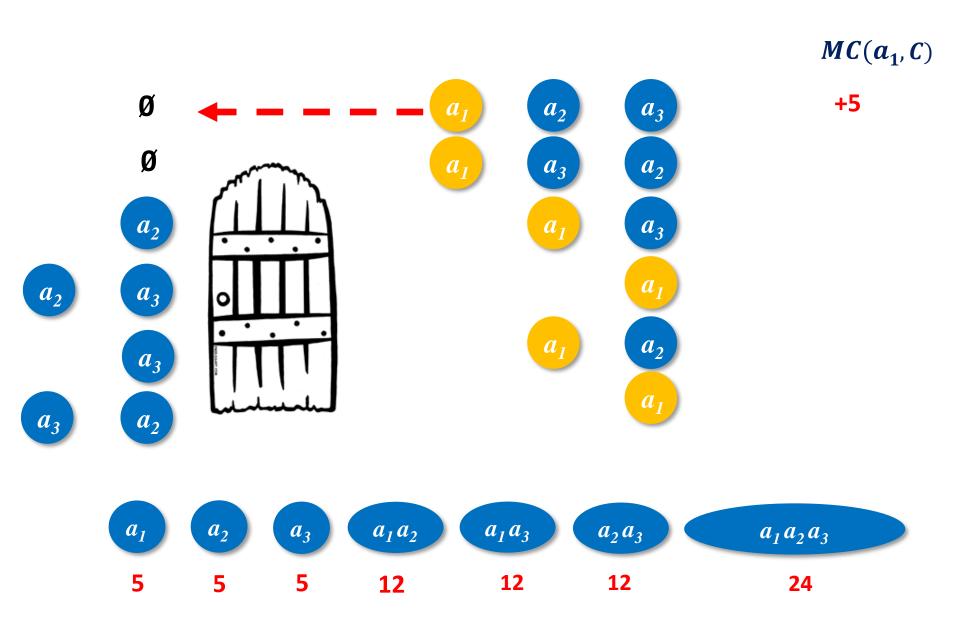


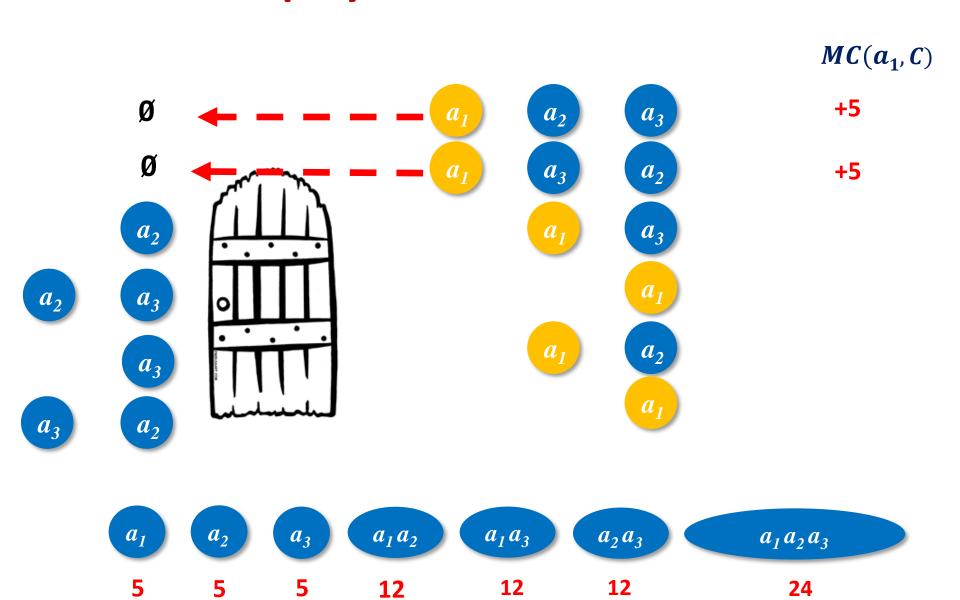


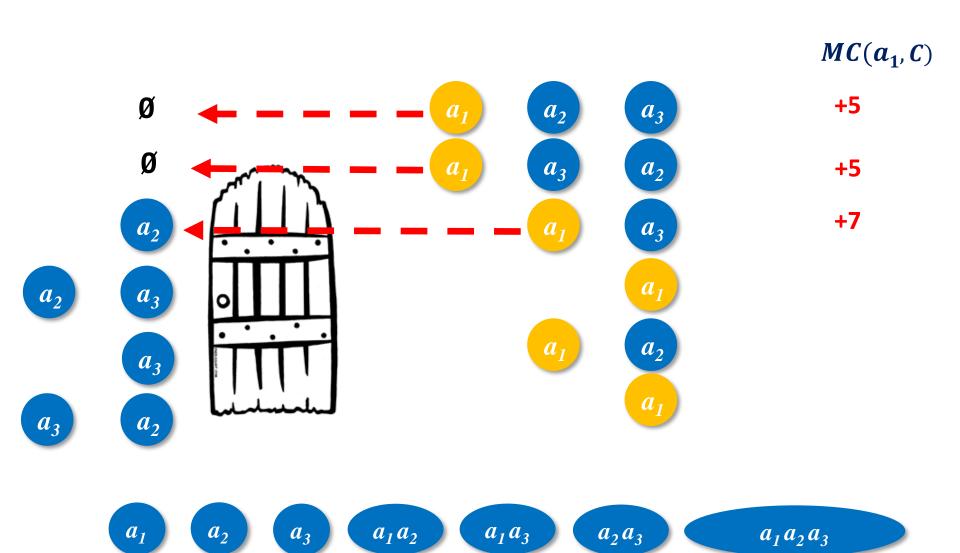


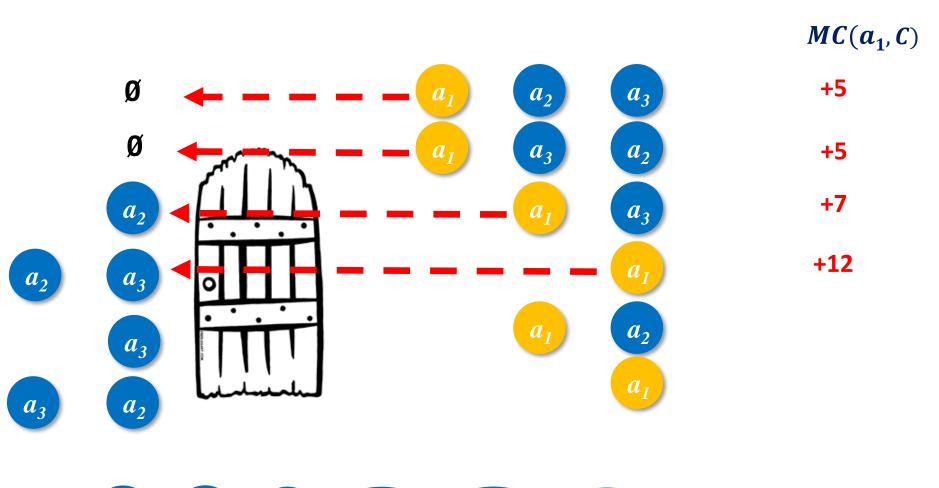




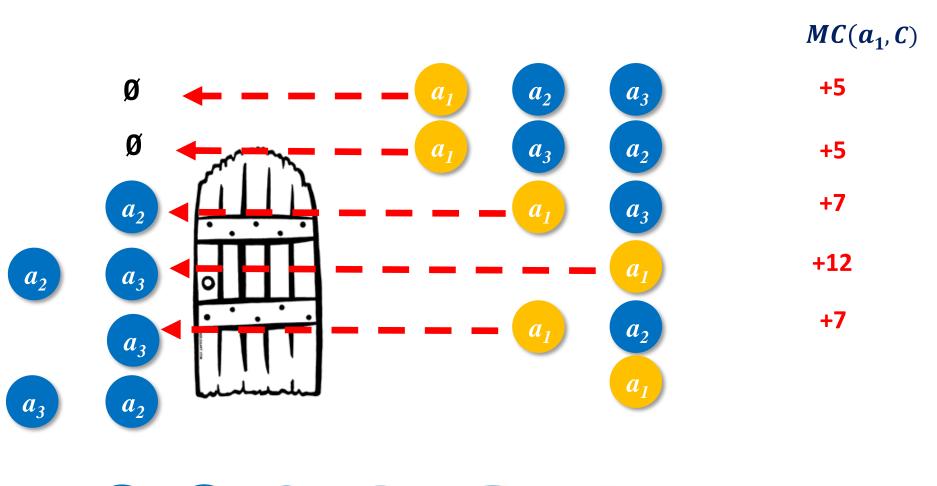




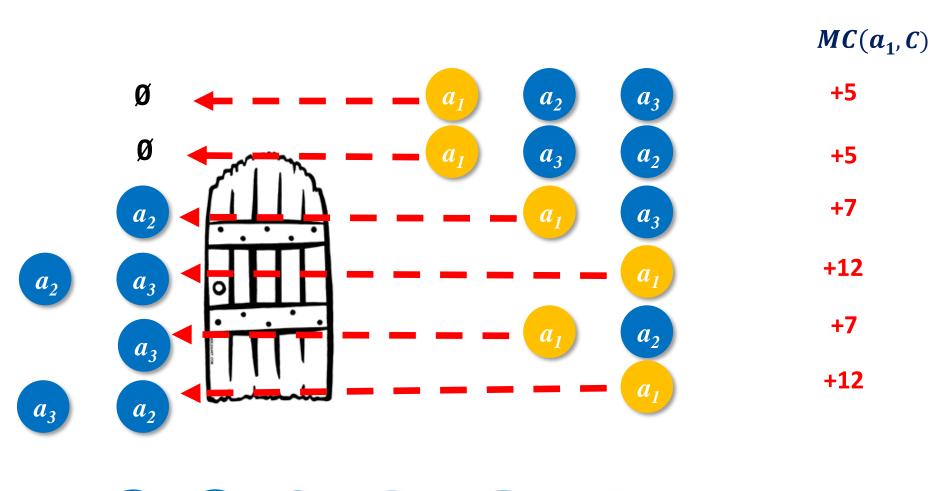




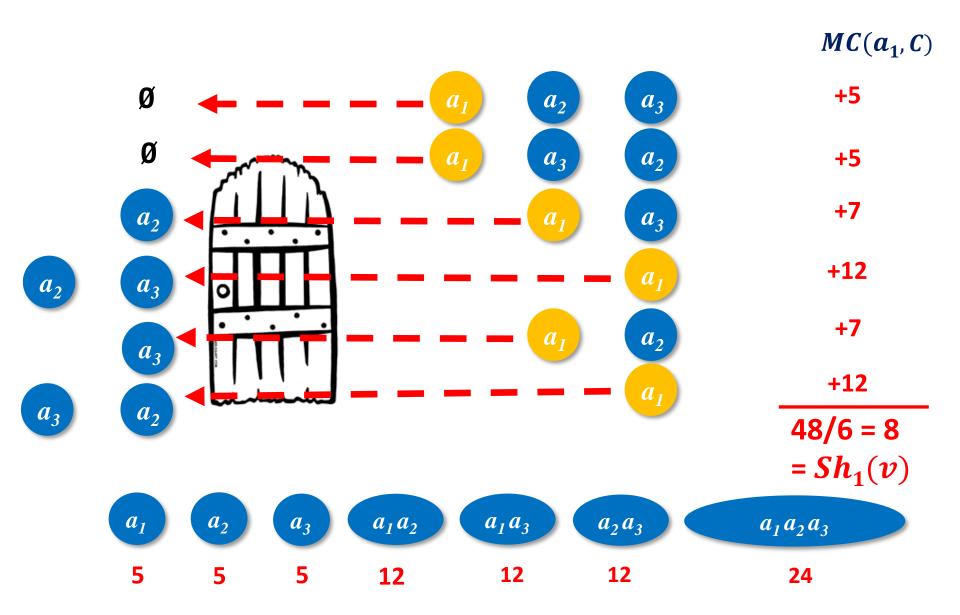


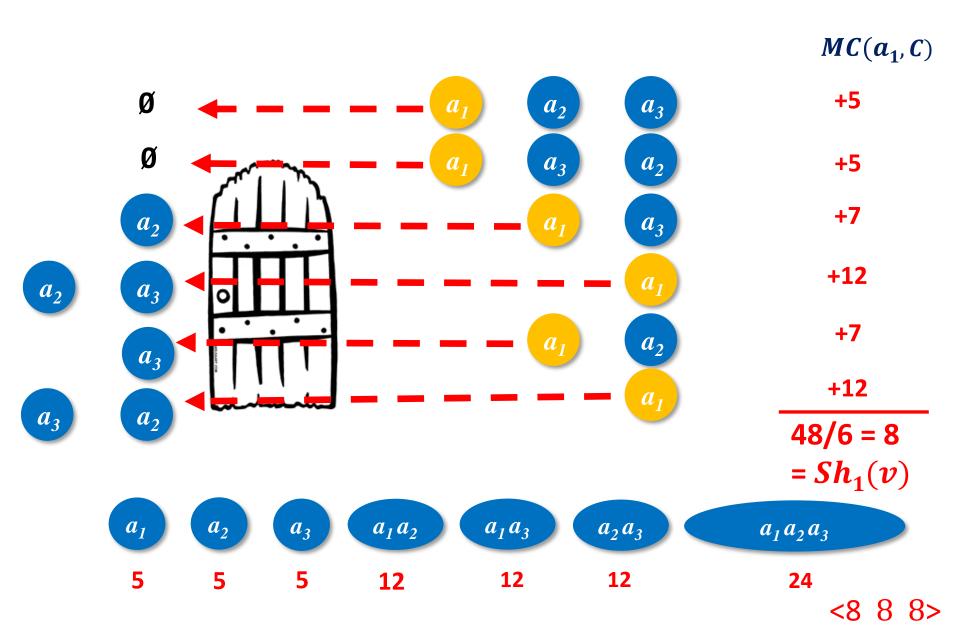












Comparison

Shapley value – weighted average marginal contribution of a player to all coalitions

$$2^{|A|} \quad Sh_i(v) = \sum_{C \subseteq A \setminus \{a_i\}} \frac{|C|! (|A| - |C| - 1)!}{|A|!} [v(C \cup \{a_i\}) - v(C)]$$

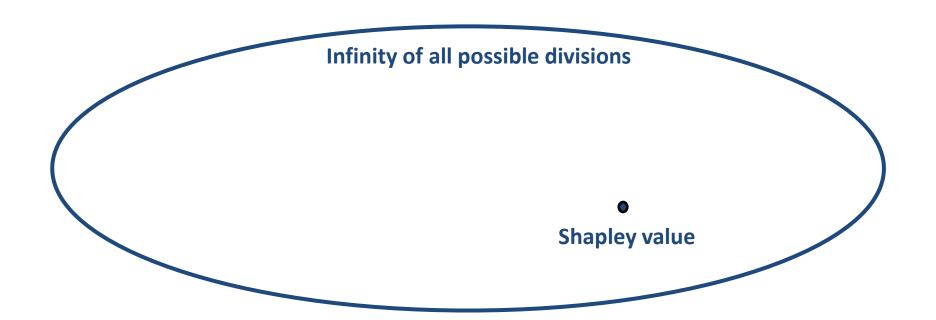
Symmetry, null player, additivity, efficiency

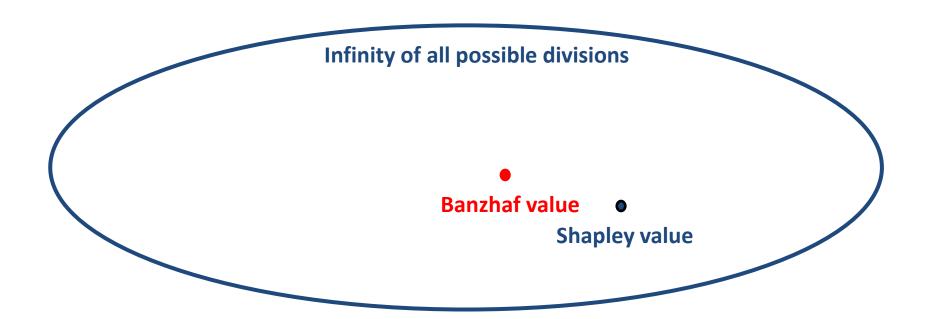
simple average marginal contribution of a player to all coalitions

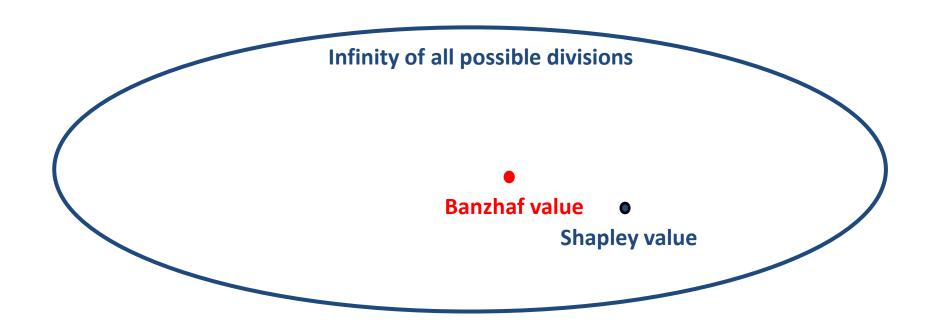
$$2^{|A|} \qquad Bh_i(v) = \frac{1}{2^{|A|-1}} \sum_{C \subseteq A \setminus \{a_i\}} (v(C \cup \{a_i\}) - v(C))$$

Symmetry, null player, additivity

Taxonomy of Solutions

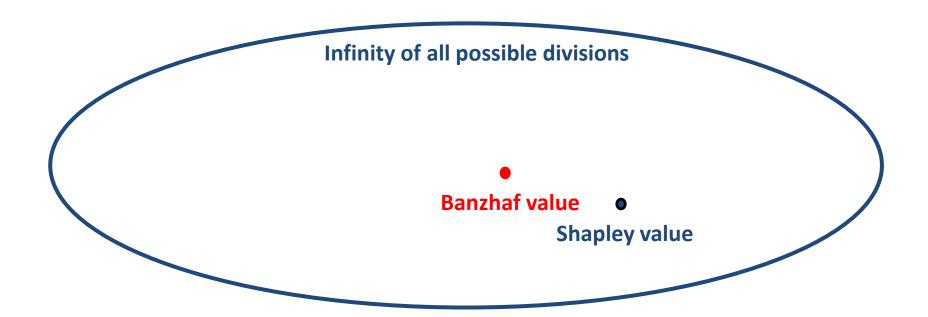




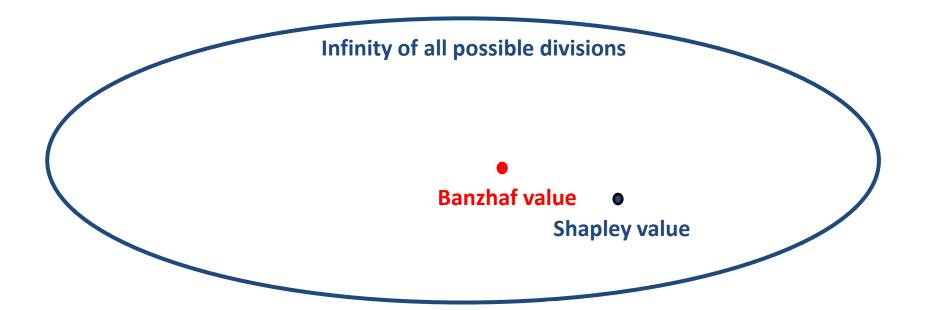


Haller (1994)

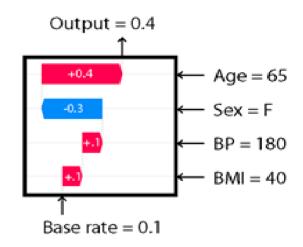
Linearity
Symmetry
Dummy player
Proxy agreement



But what about efficiency?



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Normalized Banzhaf value

$$Bh_i^N(v) = \frac{Bh_i}{\sum_{j \in A} Bh_j} v(A)$$

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van den Brink & van der Laan (1998)

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Do these axioms sound strange?

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From the literature on weighted voting games we learn that:

Even axioms that seem to be the most basic ones can lead to paradoxes

Why are supposed to weight the contribution of each feature with the total number of orderings of the present as well as the absent features?

It is not obvious why feature vector (man; 40) is different from (40;man), and why this should matter for feature importance.

From the literature on weighted voting games we learn that:

- Even axioms that seem to be the most basic ones can lead to paradoxes
- ➤ Thus, some authors recommend to use probabilistic approach when choosing between Shapley and Banzhaf value:
 - use Shapley when the order matters
 - use Banzhaf when it does not

Shapley & Banzhaf Values – Computational Challenge

$$|A|!$$
 $Sh_i(v) = \frac{1}{|A|!} \sum_{\text{all } \pi} [v(C_{\pi}(i) \cup \{a_i\}) - v(C_{\pi}(i))]$

$$2^{|A|} \quad Sh_i(v) = \sum_{\mathbf{C} \subseteq A \setminus \{a_i\}} \frac{|C|! (|A| - |C| - 1)!}{|A|!} [v(\mathbf{C} \cup \{a_i\}) - v(\mathbf{C})]$$

$$2^{|A|} \qquad Bh_i(v) = \frac{1}{2^{|A|-1}} \sum_{C \subseteq A \setminus \{a_i\}} (v(C \cup \{a_i\}) - v(C))$$





Plan of the Talk

- 1. Values in Cooperative Game Theory
- 2. Our algorithm for the Banzhaf value vs. TreeSHAP
- 3. Advantages of the Banzhaf value for tree models experimental analysis

Our Key Algorithmic Result

Shapley value

Banzhaf value

TreeSHAP Lundberg et al.

$$O(TLD^2 + n)$$



Our algorithms

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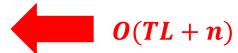
TreeSHAP Lundberg et al.

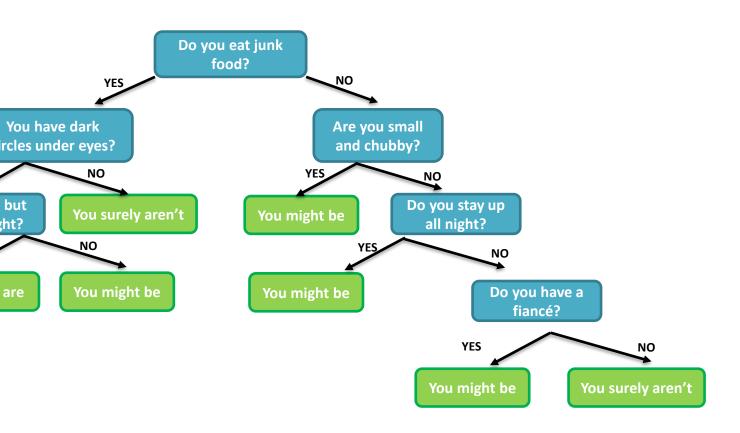
$$O(TLD^2 + n)$$



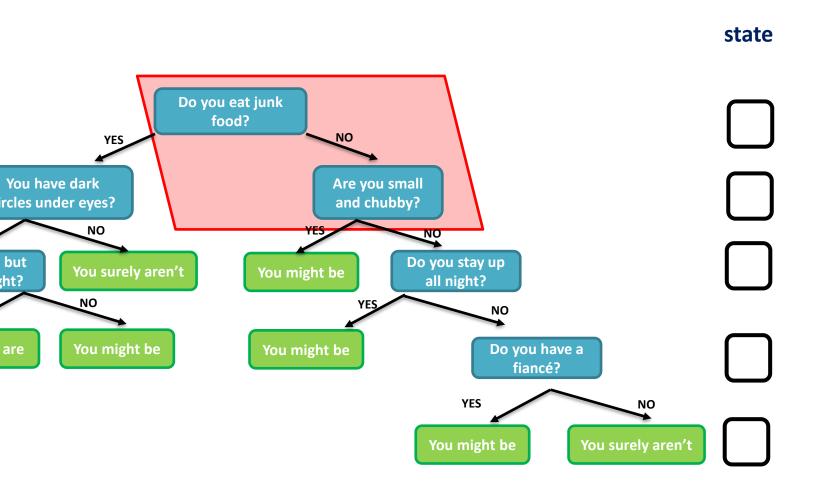
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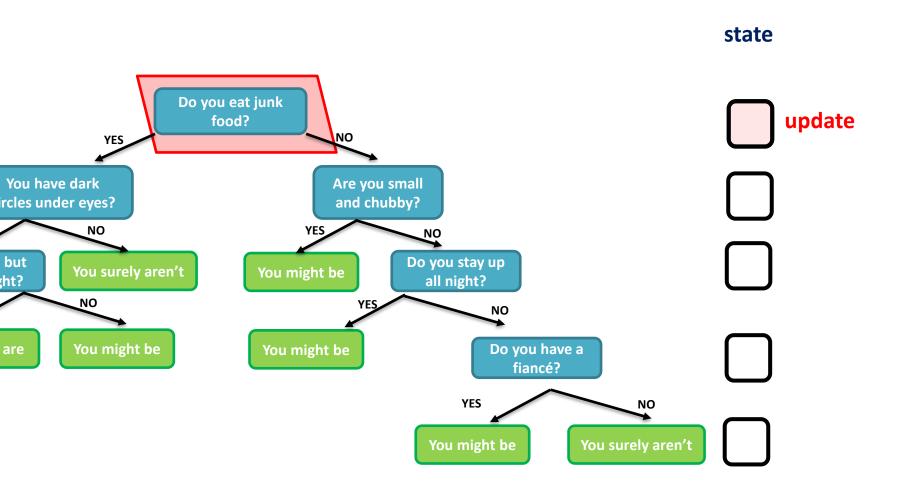




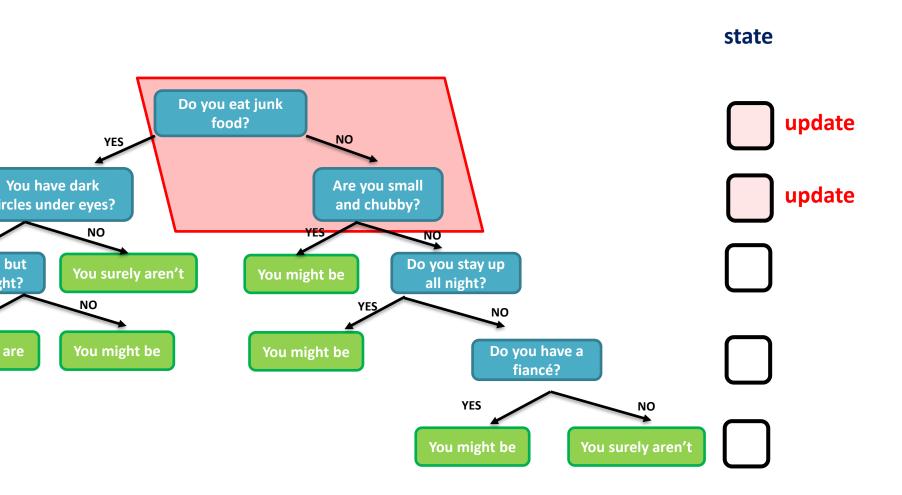
TreeSHAP is a dynamic programming algorithm.



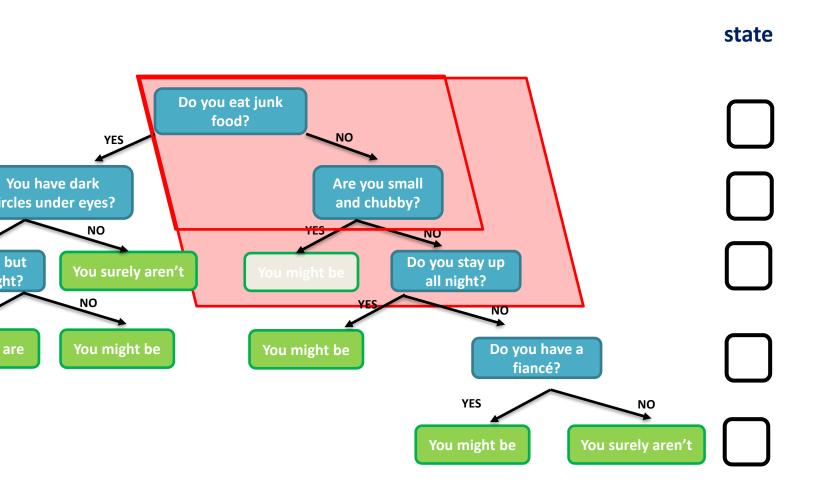
It starts from the root and goes down to the leaves extending the size of the subproblems. With each move it updates some state of up to **D** values. The update is related to the



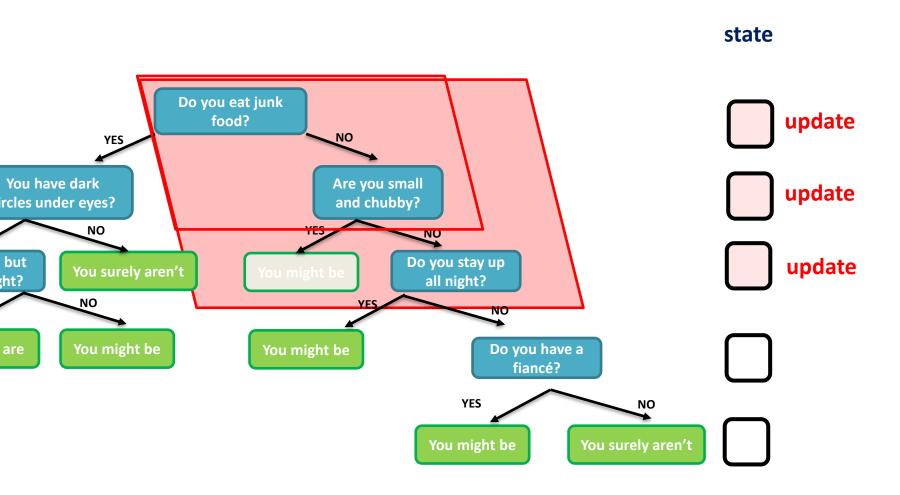
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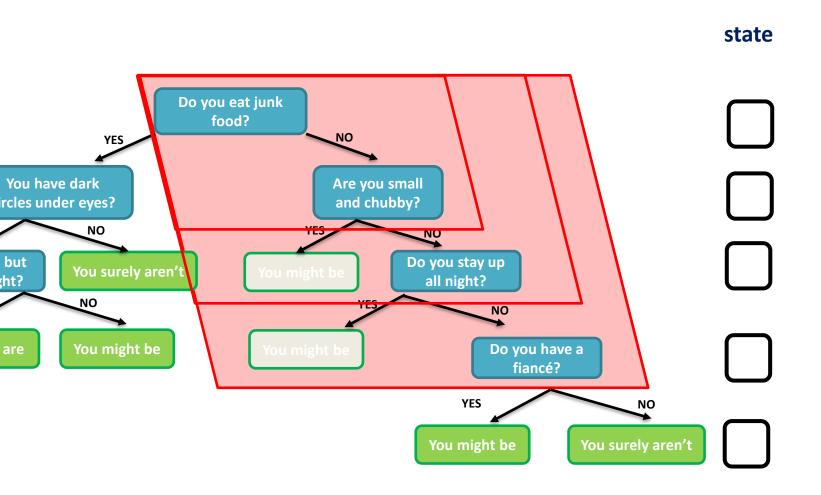
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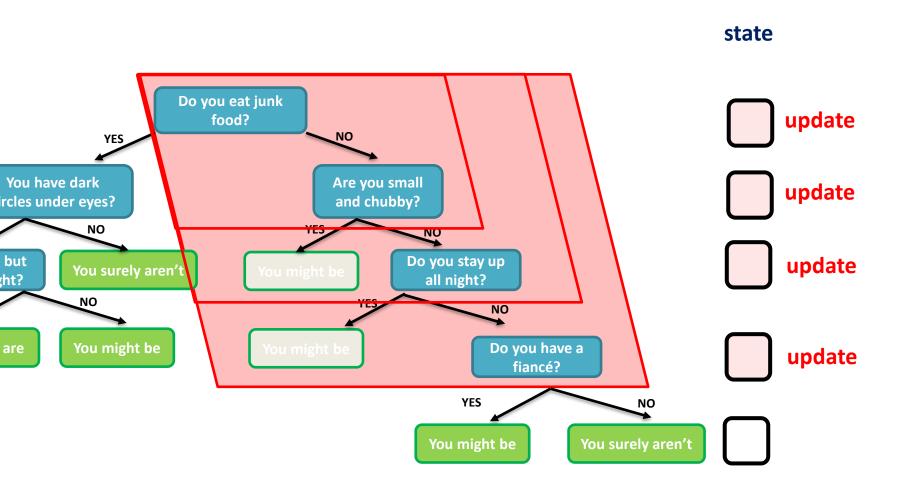
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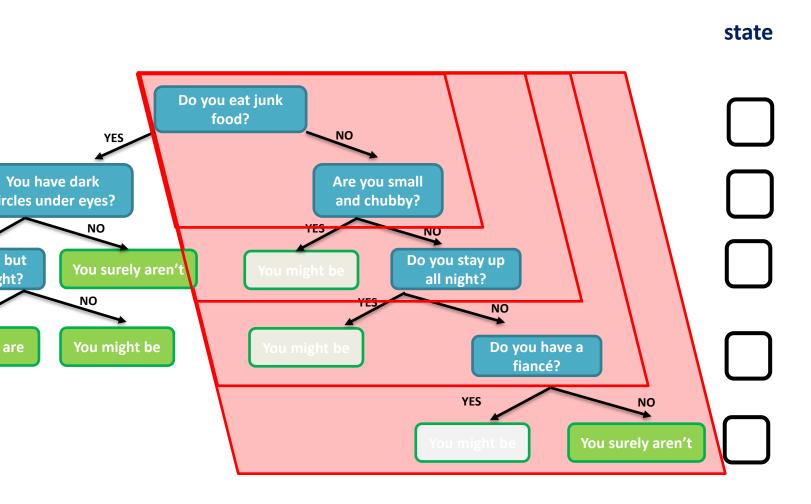
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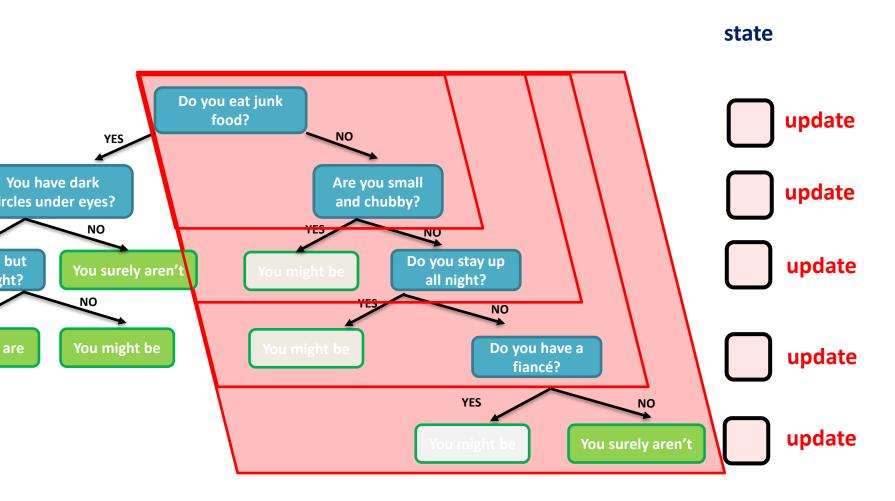
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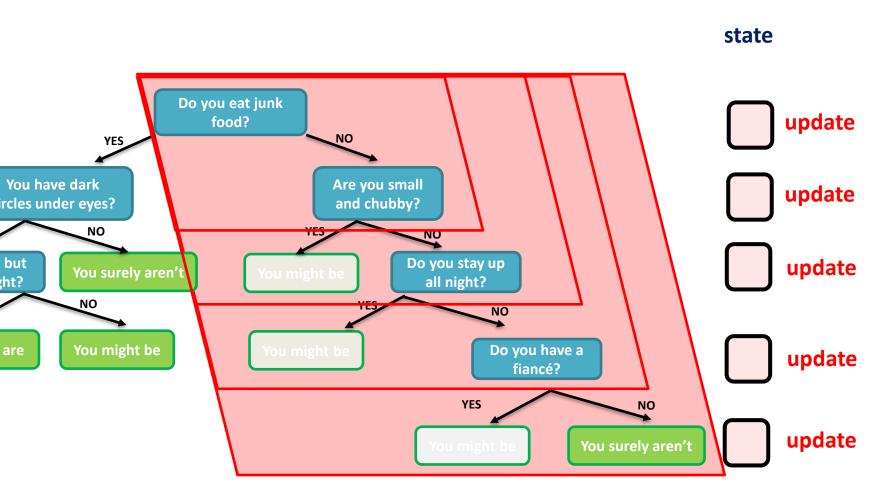
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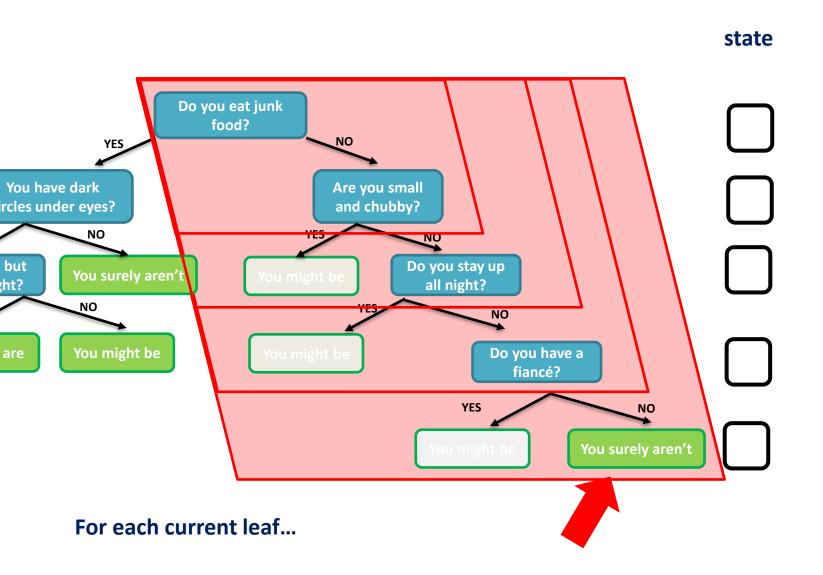
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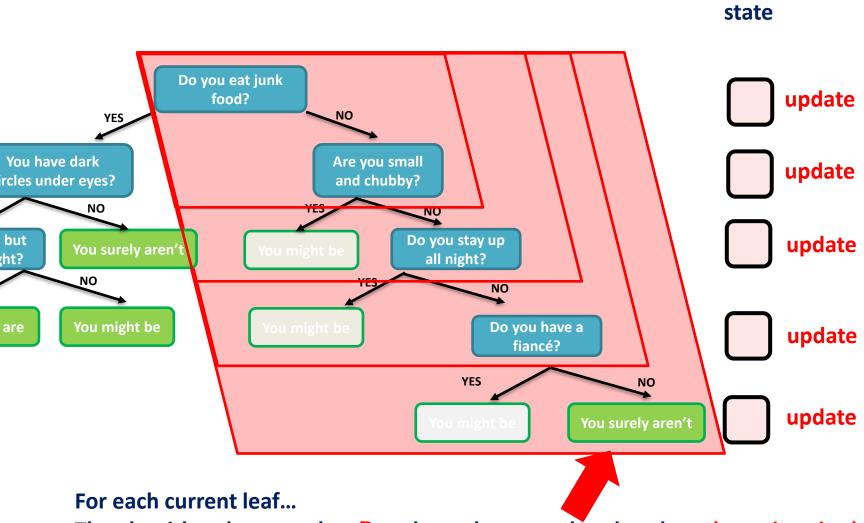


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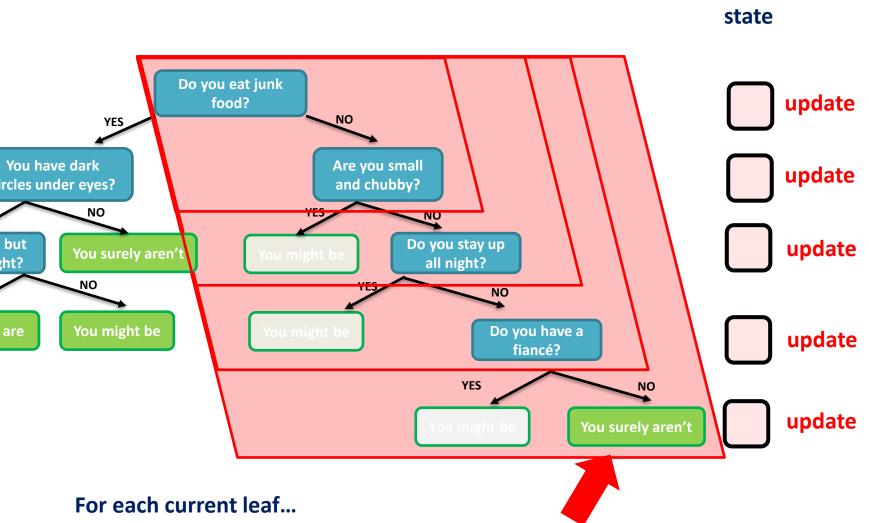


The complexity so far is O(TLD).



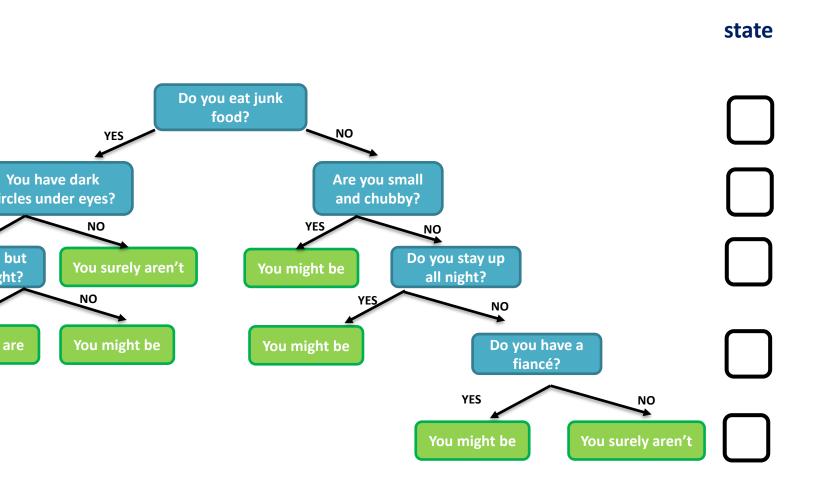


The algorithm does another **D** updates that are related to the subtractions in the marginal contributions.

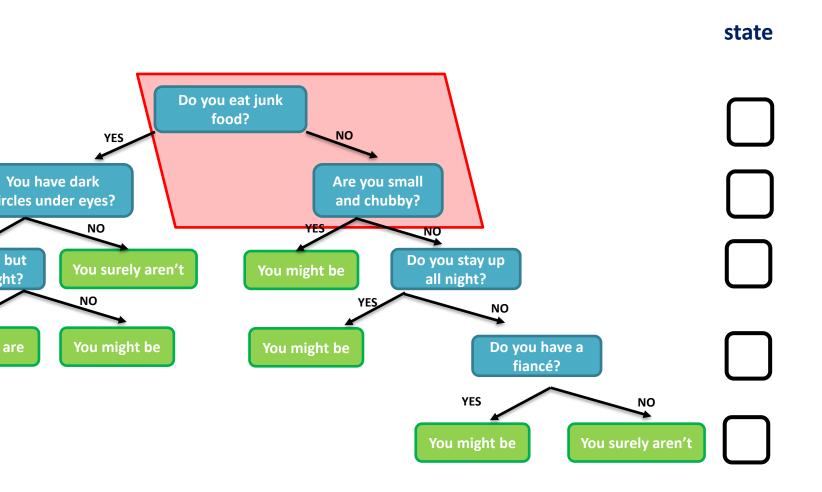


The algorithm does another D updates that are related to the subtractions in the marginal contributions. The complexity is $O(TLD^2)$.

Intuition behind our Algorithm

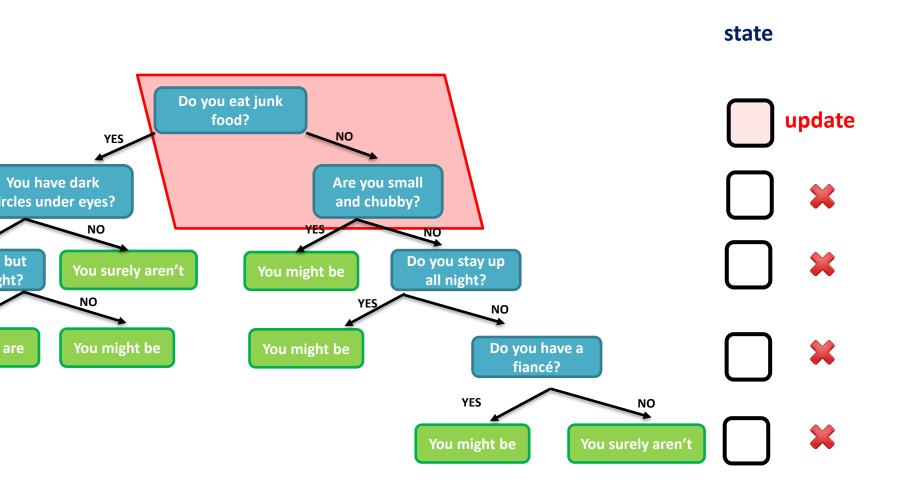


Intuition behind our Algorithm



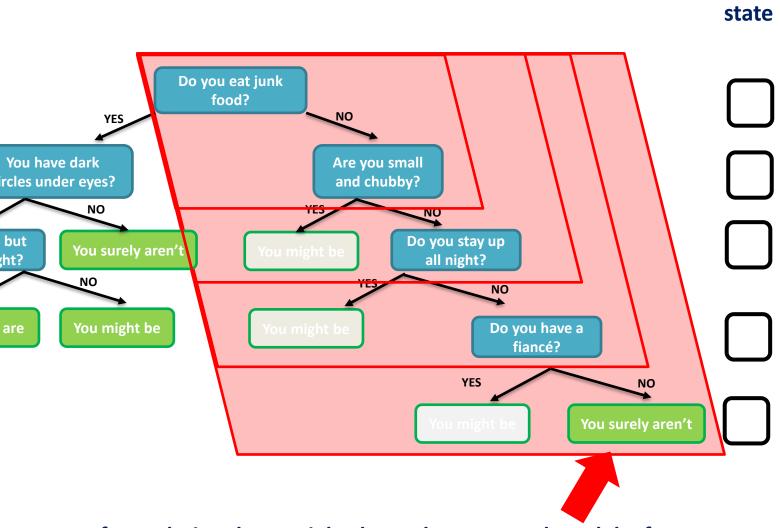
Since we don't have the weight coefficient $\frac{|C|!(|A|-|C|-1)!}{|A|!}$ in the Banzhaf value, we only update a single value for each node.

Intuition behind our Algorithm



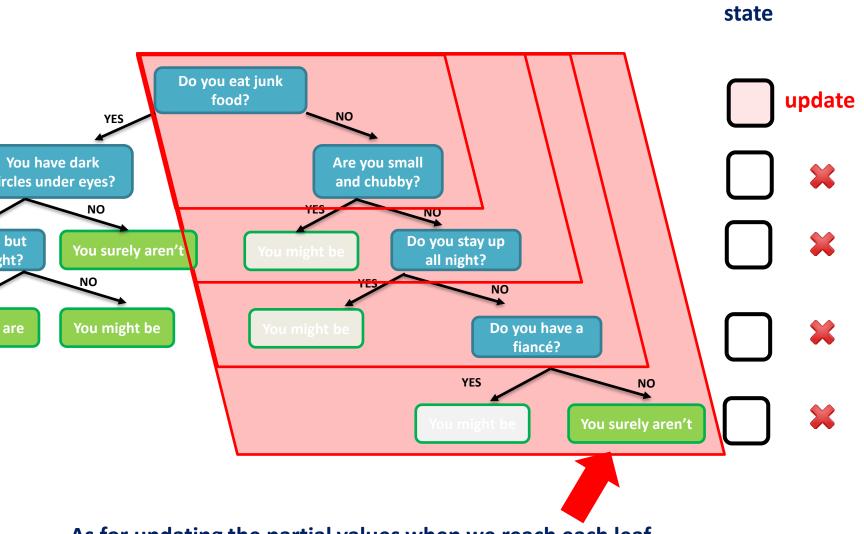
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Here the improvement comes from the definition of the Banzhaf value.



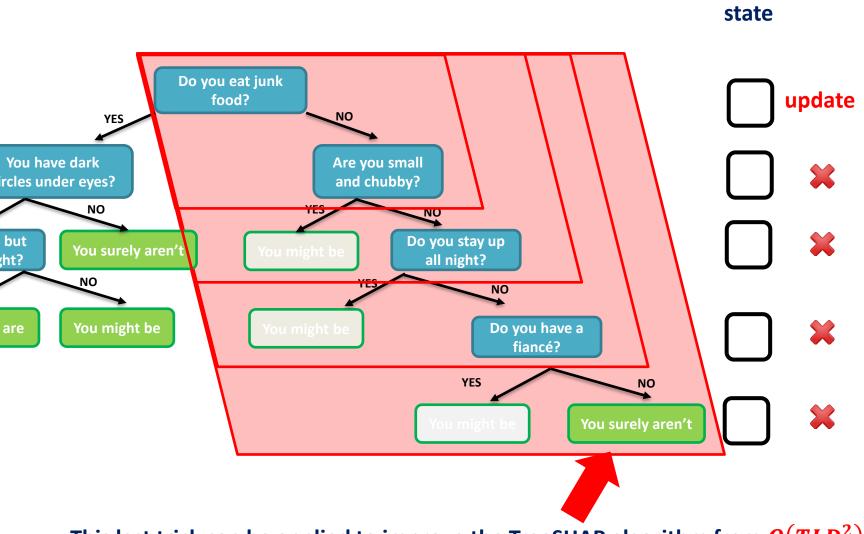
As for updating the partial values when we reach each leaf. Here, we figured out a method to do only a single update per leaf on average when we backtrack.

Intuition behind TreeSHAP



As for updating the partial values when we reach each leaf. Here, we figured out a method to do only a single update per leaf on average when we backtrack. Our complexity is O(TL).

Intuition behind TreeSHAP



This last trick can be applied to improve the TreeSHAP algorithm from $O(TLD^2)$ to O(TLD).

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Algorithms for DT & Datasets

Two arguably most popular algorithms for generating decision trees:

- sklearn implementation of Decision Trees (DT)
- xgboost implementation of Gradient Boosting Decision Trees (GB)

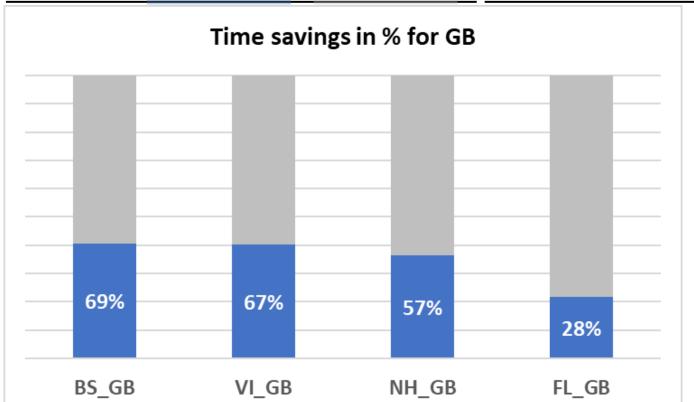
Name	Description			Approx. size				
BOSTON (BS)	This small pred information co Boston Massac price of the ho	506 rows, 13 features						
NHANES (NH)	One of the most widely-used datasets describing the health and socioeconomic status of people residing in the US.							
HEALTH_INSURANCE (HI)	A medium size dataset for predicting who might be interested in health insurance purchase. 304887 rows, 14 features							
FLIGHTS (FL)	A large dataset	for predicting	the flights' delays	. 1543718 rows, 647 features				
Dataset		xgboost		Decision tree				
	Iterations	Max depth	Learning rate	tree_depth				
BOSTON (BS)	100	6	0.01	10				
NHANES (NH)	250	4	0.2	40				
HEALTH_INS. (HI)	250 4 0.2 60			60				
FLIGHTS (FL)	250	10	0.2	100				

Experimental Results: Running Times

	BANZHAF	TREESHAP		BANZHAF	TREESHAP
BS_GB	0.48 s	0.70 s	BS_DT	0.41 s	0.41 s
VI_GB	23.63 s	35.32 s	NH_DT	3.57 s	42.87 s
NH_GB	50.20 s	1 m 28 s	VI_DT	4 m 55 s	30 m 55 s
FL_GB	13 m 18 s	48 m 8 s	FL_DT	14 m 28 s	5 h 9 m

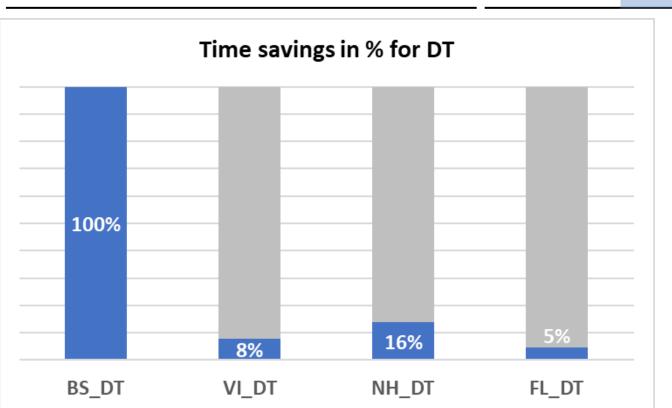
Experimental Results: Running Times

	BANZHAF	TREESHAP		BANZHAF	TREESHAP
BS_GB	0.48 s	0.70 s	BS_DT	0.41 s	0.41 s
VI_GB	23.63 s	35.32 s	NH_DT	3.57 s	42.87 s
NH_GB	50.20 s	1 m 28 s	VI_DT	4 m 55 s	30 m 55 s
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Global impact: Qualitative Difference?

Global impact – we use the same measure of global impact as Lundberg et al. (2020).

 \mathcal{D} - a dataset. $i \in A$ - feature

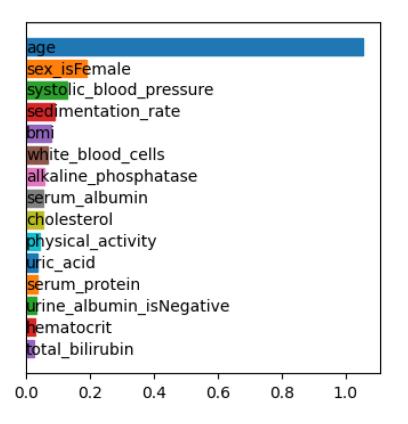
Shapley global impact:

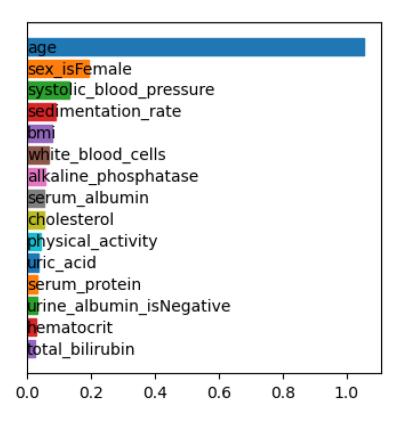
$$\Gamma_i^{\rm Sh} = \sum_{x \in \mathcal{D}} |Sh_i(x)|$$

Banzhaf global impact:

$$\Gamma_i^{\rm Sh} = \sum_{x \in \mathcal{D}} |Bh_i(x)|$$

Global Impacts: NH_GB





- (a) Global Shapley impact obtained with TREESHAP_PATH.
- (b) Global Banzhaf impact obtained with BANZHAF.

The same ordering.

Banzhaf results are virtually indistinguishable from the Shapley results. The same holds for BS and VI datasets, both GB and DT.

Global Impacts: FL_GB

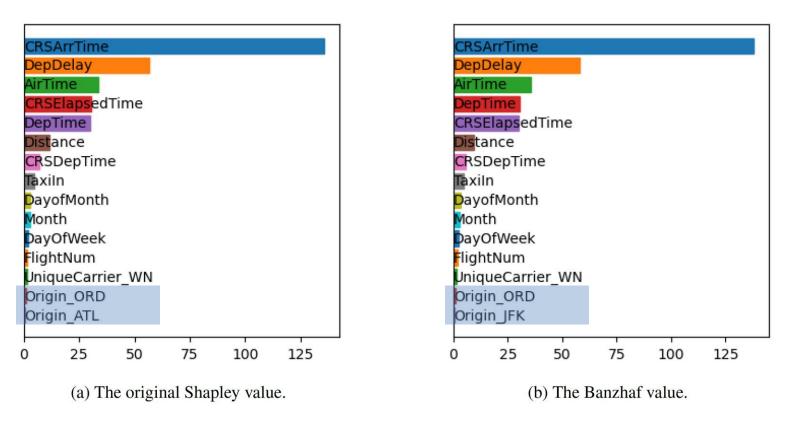


Figure 6: The global impacts of the individual features for the FLIGHTS_GB dataset. We observe small differences in the ordering.

For the largest dataset with the deepest tree, only very small differences in the ordering of features by importance can be observed for both, both GB and DT.

Individual Data Points: Qualitative Difference?

Cayley distance – the numer of swaps that are needed to generate one permutation from another

Pattern vector Disorder vector

Step 1



Cayley distance – the numer of swaps that are needed to generate one permutation from another

Pattern

Disorder

Step 1



Cayley distance – the numer of swaps that are needed to generate one permutation from another

Pattern vector 1 2 3

Disorder vector 2 5 3

 Disorder vector
 2
 5
 3
 1
 4
 6

 Select
 2
 5
 3
 1
 4
 6

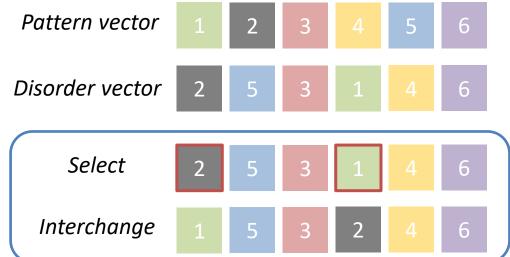
 Step 1
 Interchange
 2
 5
 3
 1
 4
 6

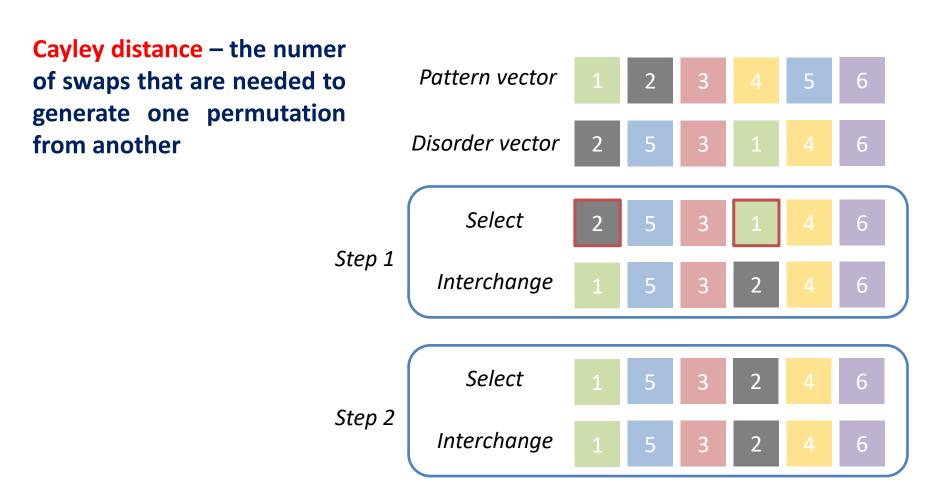
Cayley distance – the numer of swaps that are needed to generate one permutation from another

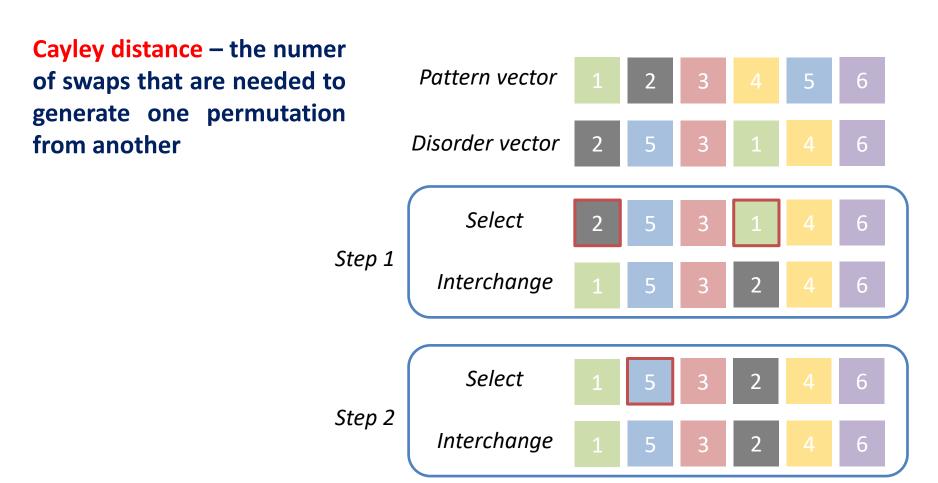
Patternation

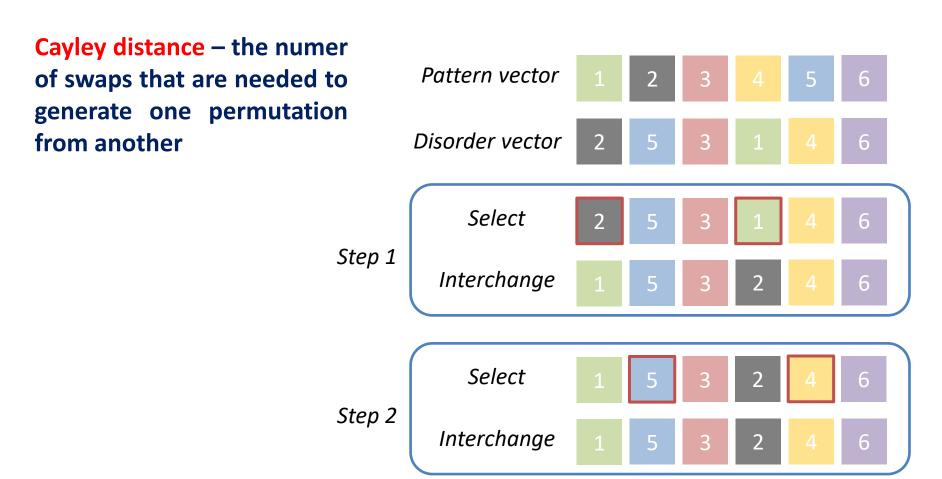
Disorder

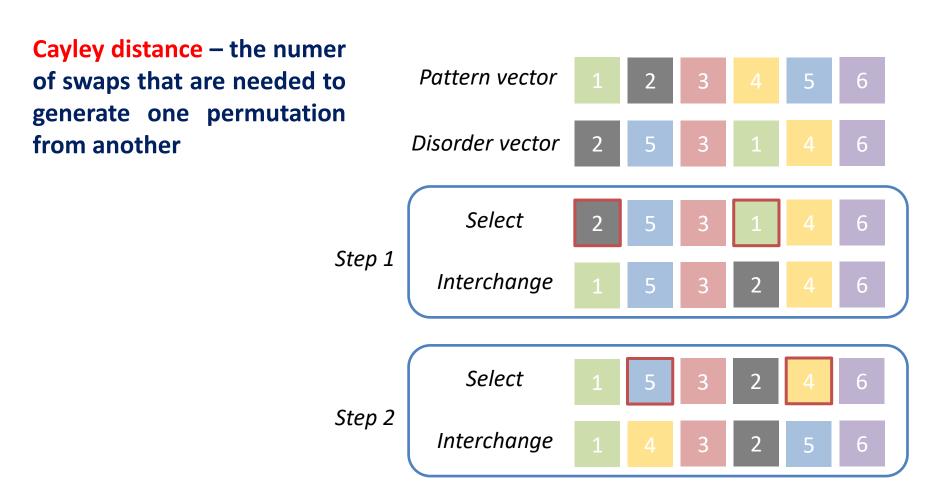
Step 1

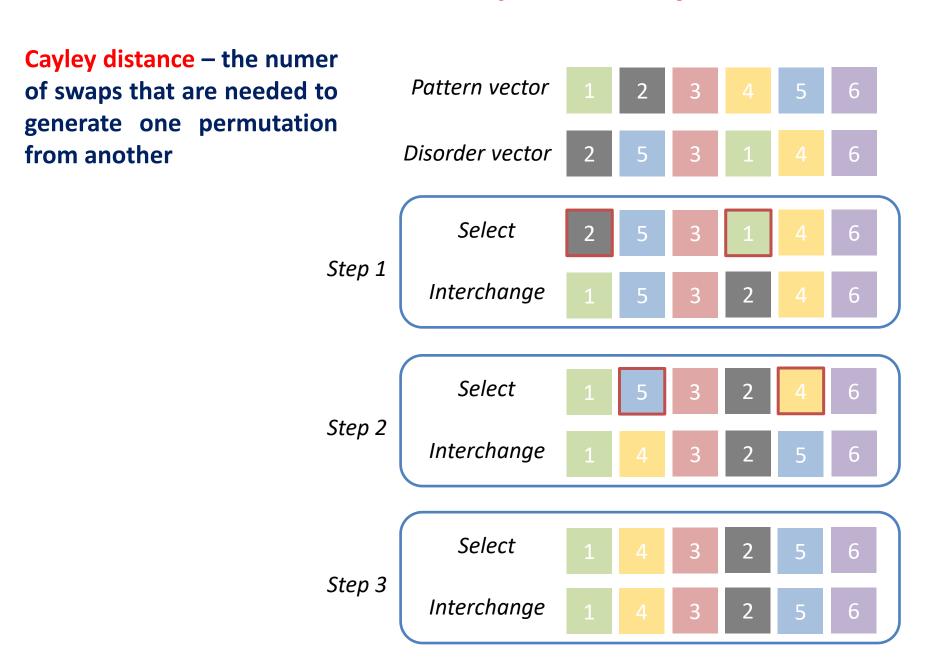


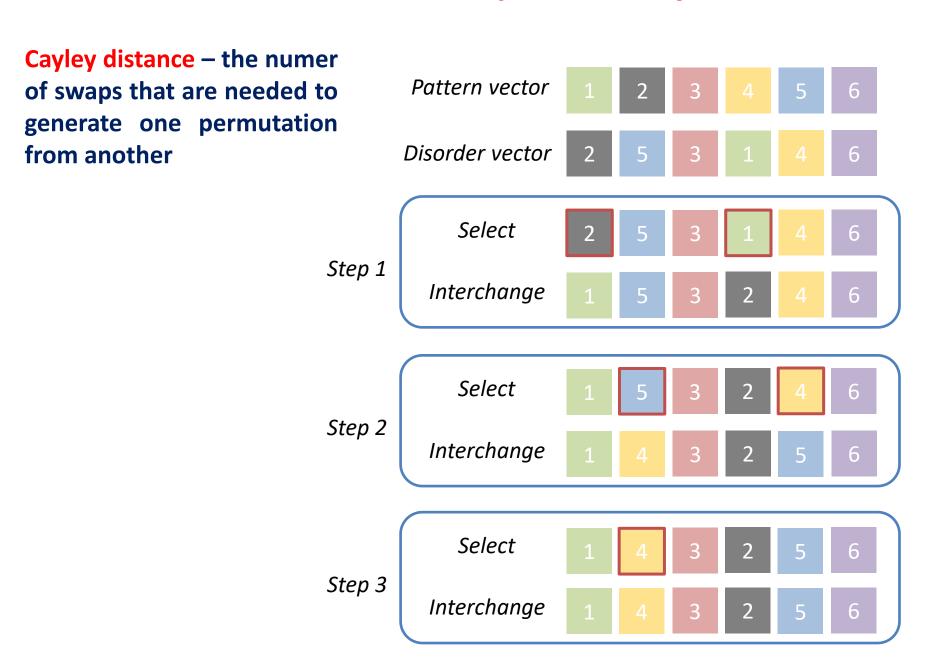


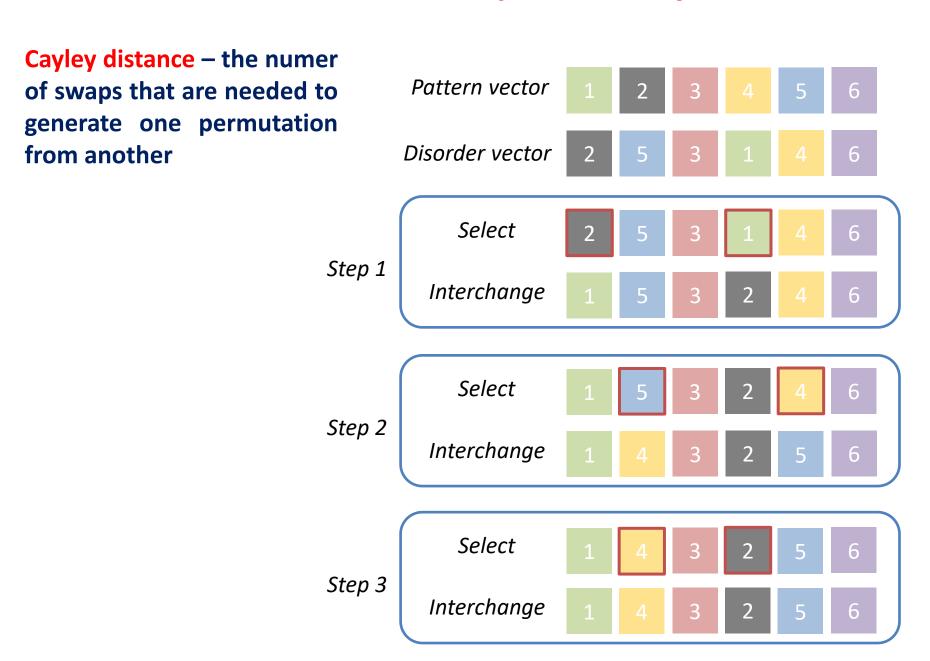


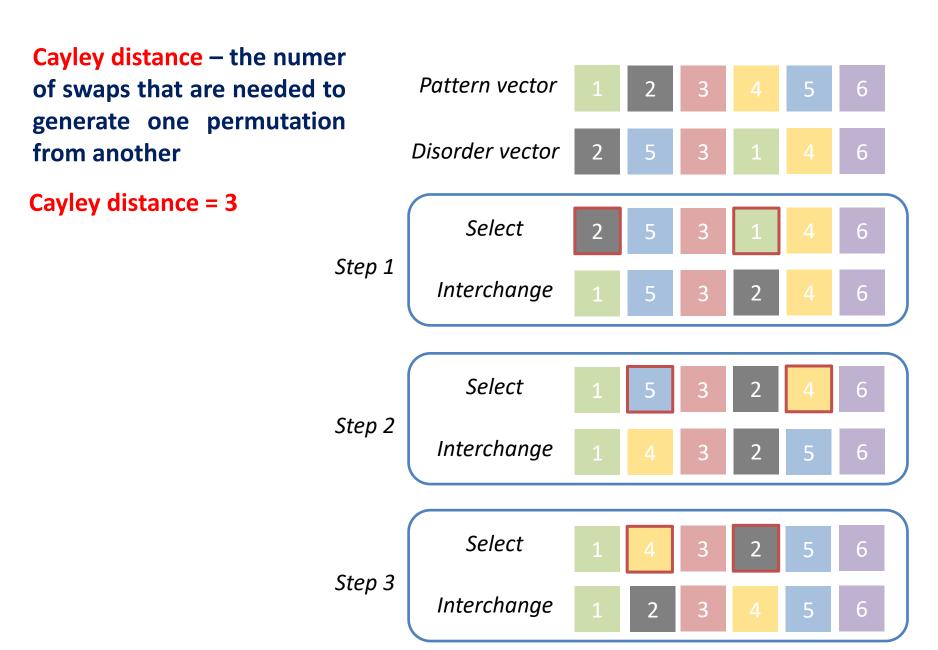


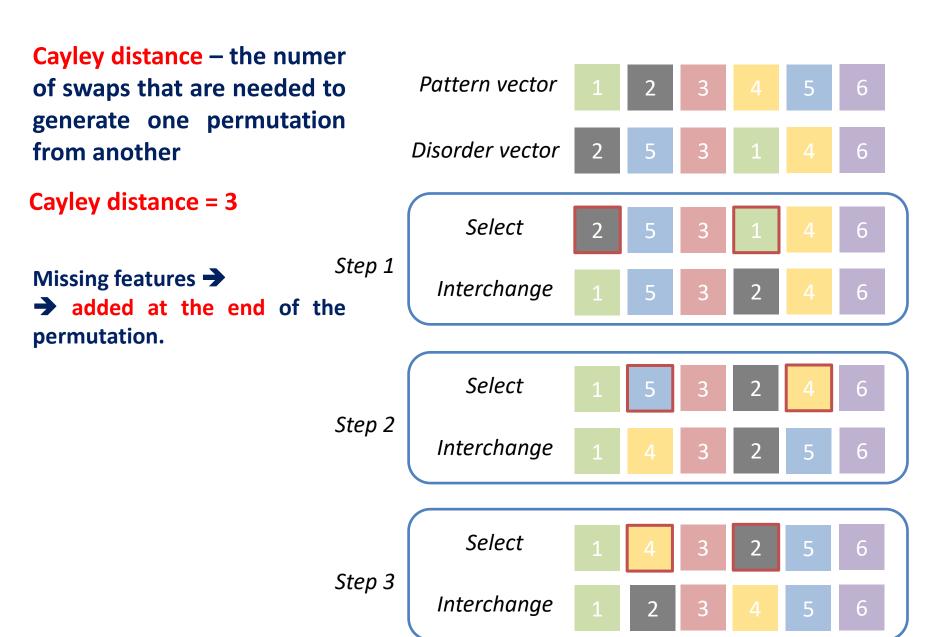












Ins/n	3	10	20	Ins/n	3	10	20
BOS_GB	0.02	1.05		BOS_DT	0.08	1.7	
NH_GB	0.01	0.34	1.53	NH_DT	0.29	3.69	10.79
VI_GB	0.02	0.73		VI_DT	0.13	2.60	
FL_GB	0.4	3.08	8.63	FL_DT	0.18	3.38	10.59

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For 98% of the data points, the respective 3 top features and their order matched.

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The orderings deviation was generally larger for DT instances, where larger tree depths were allowed.

We also studied per-feature average differences. We consider both:

- MAE (Mean Average Error) less than 5% for smaller and 20% for larger models (for top features)
- RMSE (Root Mean Square Error) for the large model the difference reached 50% even for top features

Numerical Accuracy

We can prove the following very pessimistic statement about Banzhaf values.

Lemma: The Banzhaf values can be computed with relative error at most $(1+\varepsilon)^{O(D)}-1$, where ε is machine epsilon and D is the tree depth.

This bound is quite pessimistic and at the same time not very large if double precision is used and the tree depth D is small enough.

Similar bound is impossible for TreeShap as it requires subtraction of intermediate values, which can lead to so-called catastrophic cancellations.

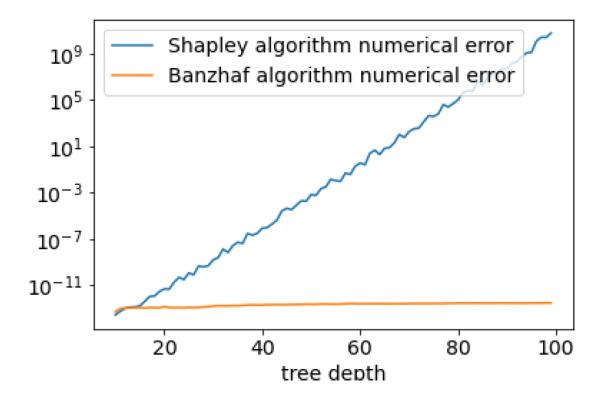
A much slower impractical algorithm for TreeShap that avoids subtractions behaves similarly to Banzhaf in experiments.

Numerical Accuracy

As more significant differences arose for large models

the algorithms might suffer numerical problems

We compare numerical stability on a simple artificially prepared instance SYNTHETIC_SPARSE for which we know the answer for both the Shapley value and the Banzhaf value.

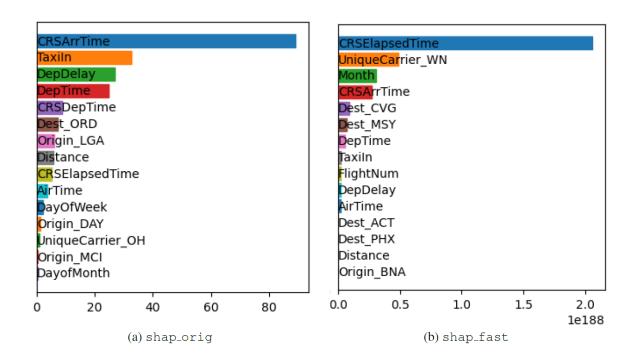


Numerical Accuracy

As more significant differences arose for large models

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When comparing different implementations of SHAP with respect to point explanations we can spot significant differences.



Take away message

Technical contribution:

We advocate the Banzhaf value for tree models:

- 1. It can be computed noticeably faster than the Shapley value
- 2. Probably more numerically stable
- 3. Both methods deliver:
 - essentially the same global impacts
 - close explanations of individual predictions

Meta level:

- Game theory and algorithmic view
- Interplay of the above areas with AI is growing
- Many more interesting problems to come

Thank you



Questions?